

Faculty of Computer Science Institute for Theoretical Computer Science, Chair for Automata Theory

# Formal Concept Analysis and Logic

## **Exercise Sheet 3**

Dr. Felix Distel Summer Semester 2012

### Exercise 9

Let  $n \ge 3$  be a natural number.

- Present a formal context such that there are n pseudo-intents, but at least 2<sup>n-1</sup> intents.
- Show that contexts of the following form have 2<sup>*n*</sup> pseudo-intents.

	<i>m</i> <sub>0</sub>	<i>m</i> <sub>1</sub>		m <sub>n</sub>	$m_{n+1}$		<i>m</i> <sub>2<i>n</i></sub>
<i>g</i> <sub>1</sub>							
			¥			¥	
<i>g</i> <sub>n</sub>							
<i>g</i> <sub>n+1</sub>	×						
:	:				$\neq$		
<b>g</b> <sub>3n</sub>	×						

Here  $\neq$  stands for the contranominal scale (where there are crosses everywhere except on the main diagonal).

#### **Exercise 10**

Complete the proof of Lemma 2.35. Show that the pseudo-closure operator  $\mathcal{D}G^*_{(G,M,I)}$  is a closure operator, i.e. it is extensive, monotone and idempotent.

Show that all intents and pseudo-intents are pseudo-closed.

## Exercise 11

Prove Lemma 2.37 which states that for a pseudo-closed set P the lectically next pseudo-closed set  $\overline{P}$  is of the form

 $P \oplus_{\mathcal{L}} x$ ,

where  $\mathcal{L} = \{ \mathcal{Q} \rightarrow \mathcal{Q}'' \mid \mathcal{Q} < P \}$ , and *x* is maximal with the property that

 $P <_{x} P \oplus_{\mathcal{L}} x.$ 

## Exercise 12

#### (optional exercise)

A set  $Q \subseteq M$  is called a *quasi-intent* iff every  $R \subseteq Q$  satisfies  $R'' \subseteq Q$  or R'' = Q''.

Prove that a set  $P \subseteq M$  is a pseudo-intent iff

- $P \neq P''$ ,
- P is a quasi-intent, and
- for all quasi-intents Q,  $Q \subsetneq P$  implies  $Q'' \subsetneq P$ .