Extracting $\mathcal{ALEQR}^{\mathsf{Self}}$-Knowledge Bases from Graphs

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Abstract. A description graph is a directed graph that has labeled vertices and edges. This document proposes a method for extracting a knowledge base from a description graph. The technique is presented for the description logic $\mathcal{ALEQR}^{\mathsf{Self}}$ which allows for conjunctions, primitive negation, existential restrictions, value restrictions, qualified number restrictions, existential self restrictions, and complex role inclusion axioms, but also sublogics may be chosen to express the axioms in the knowledge base. The extracted knowledge base entails all statements that can be expressed in the chosen description logic and are encoded in the input graph.

Keywords: description logics, formal concept analysis, terminological learning, knowledge base, general concept inclusion, canonical base, most-specific concept description, interpretation, description graph, folksonomies, social network

1 Introduction

There have been several approaches towards the combination of description logics and formal concept analysis for knowledge acquisition, knowledge exploration and knowledge completion. In [27] Rudolph invented a method for the exploration of concept inclusions holding in an $\mathcal{FLE}$-interpretation. Sertkaya provided in [28] a technique for completion of knowledge bases. Finally Distel [13] and Borchmann [8] gave a method for computing a finite base of all concept inclusions holding in an $\mathcal{EL}$-interpretation by means of formal and partial implications and the corresponding Duquenne-Guiges-base and Luxenburger-base, respectivelly.

In the following we provide a method to compute a knowledge base for concept and role inclusions holding in an $\mathcal{ALEQR}^{\mathsf{Self}}$-interpretation or description graph, respectivelly, which entails all knowledge that is encoded in the interpretation/graph and can be expressed in $\mathcal{ALEQR}^{\mathsf{Self}}$. For this purpose we need the term of a model-based most-specific concept description. Simply speaking, a most-specific concept description is a concept description, which describes a given individual x, i.e. the individual is an instance of the concept, and is most specific, i.e. for all concept descriptions C that have x as an individual, the most specific concept is subsumed by C. Since we do not want to use greatest fixpoint semantics here, we restrict the role-depth to ensure existence of most specific concepts. We chose $\mathcal{ALEQR}^{\mathsf{Self}}$ since it is an expressive description logic that does not allow for disjunctions (like $\mathcal{ALC}$) and hence will not model the examples in the input graph too exactly.

We start with a short introduction on the description logic $\mathcal{ALEQR}^{\mathsf{Self}}$. Then we define description graphs and show their equivalence to interpretations. Furthermore we then present model-based most-specific concept descriptions and their relationships to formal concept analysis. Please note that many of the results on model-based most-specific concept descriptions and bases of concept inclusions have already been observed and proven by Distel [13] and Borchmann [8] for the description logic $\mathcal{EL}^{\perp}$, which allows the bottom concept, concept conjunction and existential restriction. We then continue with induced concept and role contexts and eventually construct the knowledge base. Therefore we extend the previous results to the additional concept constructors of $\mathcal{ALEQR}^{\mathsf{Self}}$ and also take complex role inclusions into account.

2 The Description Logic $\mathcal{ALEQR}^{\mathsf{Self}}$

Let (N_C, N_R) be a *signature*, i.e. N_C is a set of *concept names* and N_R is a set of *role names*, such that N_C and N_R are disjoint. We stick to the usual notations and hence concept names are written as upper-case latin letters, e.g. A and B, and role names are written as lower-case latin letters, e.g. r and s. An *interpretation* over (N_C, N_R) is a tuple $\mathcal{I} = (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$, where $\Delta^{\mathcal{I}}$ is a non-empty set, called *domain*, and $\cdot^{\mathcal{I}}$ is an *extension function* that maps concept names $A \in N_C$ to subsets $A^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}}$ and role names $r \in N_R$ to binary relations $r^{\mathcal{I}} \subseteq \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}}$.

The set of all $\mathcal{ALEQR}^{\mathsf{Self}}$-*concept descriptions* is denoted by $\mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R)$ and is inductivelly defined as follows. Every concept name $A \in N_C$ and $\perp$, $\top$ are *atomic* $\mathcal{ALEQR}^{\mathsf{Self}}$-concept descriptions. If $A \in N_C$ is a concept name, $r \in N_R$ is a role name, $C, D \in \mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R)$ are concept descriptions and $n \in \mathbb{N}_+$ is a positive integer, then $\neg A$, $C \sqcap D$, $\exists\, r.\, C$, $\forall\, r.\, C$, $\geq n.\, r.\, C$, $\leq n.\, r.\, C$ and $\exists\, r.\, \mathsf{Self}$ are *complex* $\mathcal{ALEQR}^{\mathsf{Self}}$-concept descriptions. The extension function of an interpretation $\mathcal{I}$ is then extended to all $\mathcal{ALEQR}^{\mathsf{Self}}$-concept descriptions as shown in the semantics column of figure 1.

name	*syntax C*	*semantics* $C^{\mathcal{I}}$
bottom concept	$\perp$	$\emptyset$
top concept	$\top$	$\Delta^{\mathcal{I}}$
primitive negation	$\neg A$	$\Delta^{\mathcal{I}} \setminus A^{\mathcal{I}}$
conjunction	$C \sqcap D$	$C^{\mathcal{I}} \cap D^{\mathcal{I}}$
existential restriction	$\exists\, r.\, C$	$\left\{ x \in \Delta^{\mathcal{I}} \mid \exists y \in \Delta^{\mathcal{I}} \colon (x,y) \in r^{\mathcal{I}} \wedge y \in C^{\mathcal{I}} \right\}$
value restriction	$\forall\, r.\, C$	$\left\{ x \in \Delta^{\mathcal{I}} \mid \forall y \in \Delta^{\mathcal{I}} \colon (x,y) \in r^{\mathcal{I}} \to y \in C^{\mathcal{I}} \right\}$
qualified number	$\geq n.\, r.\, C$	$\left\{ x \in \Delta^{\mathcal{I}} \,\middle|\, \exists Y \in \binom{\Delta^{\mathcal{I}}}{n} \colon \{x\} \times Y \subseteq r^{\mathcal{I}} \wedge Y \subseteq C^{\mathcal{I}} \right\}$
restriction	$\leq n.\, r.\, C$	$\left\{ x \in \Delta^{\mathcal{I}} \,\middle|\, \forall Y \in \binom{\Delta^{\mathcal{I}}}{n+1} \colon \{x\} \times Y \subseteq r^{\mathcal{I}} \to Y \not\subseteq C^{\mathcal{I}} \right\}$
self restriction	$\exists\, r.\, \mathsf{Self}$	$\left\{ x \in \Delta^{\mathcal{I}} \mid (x,x) \in r^{\mathcal{I}} \right\}$

Fig. 1. Concept Constructors of $\mathcal{ALEQR}^{\mathsf{Self}}$

Note that every individual, that has no r-successors in the interpretation $\mathcal{I}$ at all, is an element of the extension of every value restriction $\forall\, r.\, C$ for arbitrary concept descriptions C. We use the usual notation $\binom{X}{k}$ for the set of all subsets of X with exactly k elements, i.e. $\binom{X}{k} = \{Y \subseteq X \mid |Y| = k\} \subseteq 2^X$ holds. It is well known that $\left|\binom{X}{k}\right| = \binom{|X|}{k}$ holds.

Furthermore $\mathcal{ALEQR}^{\mathsf{Self}}$ allows to express the following *terminological axioms*. When A is a concept name and C, D are concept descriptions, then $C \sqsubseteq D$ is a *(general) concept inclusion* (abbr. *GCI*) and $A \equiv C$ is a *concept definition*. Of course every concept definition $A \equiv C$ can be simulated by two concept inclusions $A \sqsubseteq C$ and $C \sqsubseteq A$. If $r, r_1, \ldots, r_n, s$ are role names, then $r \sqsubseteq s$ is a *simple role inclusion* and $r_1 \circ \ldots \circ r_n \sqsubseteq s$ is a *complex role inclusion*, also called *role inclusion axiom* (abbr. *RIA*). We then say that an interpretation $\mathcal{I}$ is a *model* of an axiom α, denoted as $\mathcal{I} \models \alpha$, if the condition in the semantics column of figure 2 is satisfied.

name	*syntax* α	*semantics* $\mathcal{I} \models \alpha$
concept inclusion	$C \sqsubseteq D$	$C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$
concept definition	$A \equiv C$	$A^{\mathcal{I}} = C^{\mathcal{I}}$
simple role inclusion	$r \sqsubseteq s$	$r^{\mathcal{I}} \subseteq s^{\mathcal{I}}$
complex role inclusion	$r_1 \circ r_2 \circ \ldots \circ r_n \sqsubseteq s$	$r_1^{\mathcal{I}} \circ r_2^{\mathcal{I}} \circ \ldots \circ r_n^{\mathcal{I}} \subseteq s^{\mathcal{I}}$

Fig. 2. Axiom Constructors of $\mathcal{ALEQR}^{\mathsf{Self}}$

$\circ$ denotes the *product* operator where $R \circ S := \{(x, z) \mid \exists y\colon (x, y) \in R \wedge (y, z) \in S\}$.

A *TBox* is a set of concept inclusions and concept definitions, and a *RBox* is a set of role inclusions. $\mathcal{I}$ is a *model* of a TBox $\mathcal{T}$, denoted as $\mathcal{I} \models \mathcal{T}$, if $\mathcal{I}$ is a model of all axioms $\alpha \in \mathcal{T}$, and analogously for RBoxes $\mathcal{R}$. A *knowledge base* $\mathcal{K}$ is a pair $(\mathcal{T}, \mathcal{R})$ of a TBox $\mathcal{T}$ and a RBox $\mathcal{R}$.

Definition 1 (Knowledge Base). *Let $\mathcal{I}$ be an interpretation. A* knowledge base *for $\mathcal{I}$ is a knowledge base $\mathcal{K}$ that has the following properties.*

(sound) *All axioms in $\mathcal{K}$ hold in $\mathcal{I}$, i.e. $\mathcal{I} \models \mathcal{K}$.*
(complete) *All axioms that hold in $\mathcal{I}$, are entailed by $\mathcal{K}$, i.e. $\mathcal{I} \models \alpha \Rightarrow \mathcal{K} \models \alpha$.*
(irredundant) *None of the axioms in $\mathcal{K}$ follows from the others, i.e. $\mathcal{K} \setminus \{\alpha\} \not\models \alpha$ holds for all $\alpha \in \mathcal{K}$.*

3 Graphs

The semantics of $\mathcal{ALEQR}^{\mathsf{Self}}$ can also be characterized by means of description graphs, which are cryptomorphic to interpretations. A *description graph* over (N_C, N_R) is a tuple $\mathcal{G} = (V, E, \ell)$, such that the following conditions hold.

1. (V, E) is a directed graph, i.e. V is a set of *vertices* and $E \subseteq V \times V$ is a set of *directed edges* on V. For an edge $(v, w) \in E$ we say that v and w are connected, v is the *source vertex* and w is the *target vertex* of (v, w).
2. $\ell = \ell_V \,\dot\cup\, \ell_E$ is a *labeling function*, where $\ell_V\colon V \to 2^{N_C}$ maps each vertex $v \in V$ to a label set $\ell_V(v) \subseteq N_C$, and $\ell_E\colon E \to 2^{N_R}$ maps each edge $(v, w) \in E$ to a label set $\ell_E(v, w) \subseteq N_R$.

The vertices of the graph $\mathcal{G}$ are labeled with subsets of N_C, to indicate the concept names they belong to. Analogously, the edges are labeled with subsets of N_R to allow multiple relations between the same two vertices in the graph. Usually, one would also define a root note $v_0 \in V$ for description graphs, but this is not neccessary for our purposes here.

A description graph may also be called *folksonomy* or *social network* here. For example, the set N_R of role names in the signature may contain a relation friend, that connects friends in a social network (graph). Other relations are for example isMarriedWith, sentFriendrequestTo, likes, follows, and hasAttendedEvent, with their obvious meaning. The vertices in a social network are of course the users. The vertex labels in the set N_C of concept names can be used to categorize the users in a social network, e.g. by nationality, sex, marital status, and profession.

For each description graph $\mathcal{G} = (V, E, \ell)$ we can define a canonical interpretation $\mathcal{I}_\mathcal{G}$ that contains all information that is provided in $\mathcal{G}$, as follows. The domain is just the vertex set, i.e. $\Delta^{\mathcal{I}_\mathcal{G}} := V$, and the extensions of concept names $A \in N_C$ and of role names $r \in N_R$, respectivelly, are given as follows.

$$\begin{aligned} A^{\mathcal{I}_\mathcal{G}} &:= \ell_V^{-1}(A) = \{v \in V \mid A \in \ell(v)\} \\ r^{\mathcal{I}_\mathcal{G}} &:= \ell_E^{-1}(r) = \{(v, w) \in E \mid r \in \ell(v, w)\} \end{aligned}$$

Furthermore, we can easily construct a description graph $\mathcal{G}_\mathcal{I}$ from an interpretation $\mathcal{I} = (\Delta^\mathcal{I}, \cdot^\mathcal{I})$ by setting $\mathcal{G}_\mathcal{I} := (V, E, \ell)$ where

$$\begin{aligned} V &:= \Delta^\mathcal{I} & \ell_V(v) &:= \left\{ A \in N_C \,\middle|\, v \in A^\mathcal{I} \right\} \\ E &:= \bigcup_{r \in N_R} r^\mathcal{I} & \ell_E(v, w) &:= \left\{ r \in N_R \,\middle|\, (v, w) \in r^\mathcal{I} \right\}. \end{aligned}$$

It can be readily verified that both transformations are inverse to each other, i.e. $\mathcal{I}_{\mathcal{G}_\mathcal{I}} = \mathcal{I}$ holds for all interpretations $\mathcal{I}$, and $\mathcal{G}_{\mathcal{I}_\mathcal{G}} = \mathcal{G}$ holds for all description graphs $\mathcal{G}$.

As a consequence we do not have to distinguish between interpretations and description graphs, and we may also compute model-based most-specific concept descriptions (which are usually defined for individuals of an interpretation, cf. next section) for vertices in description graphs. In the following we want to propose a method to compute a knowledge base $\mathcal{K} = (\mathcal{T}, \mathcal{R})$ from a given description graph $\mathcal{G}$ that entails all knowledge that is encoded in $\mathcal{G}$ and is expressible in the description logic $\mathcal{ALEQR}^{\mathsf{Self}}$.

4 Most-Specific Concept Descriptions

The *role depth* $\mathsf{rd}(C)$ of a concept description C is defined as the greatest number of roles in a path in the syntax tree of C. Formally we inductivelly define the role depth as follows.

1. Every simple concept description $A, \bot, \top$ and every primitive negation $\neg A$ has role depth 0.
2. The role depth of a conjunction is the maximum of the role depths of the conjuncts, i.e. $\mathsf{rd}(C \sqcap D) := \mathsf{rd}(C) \vee \mathsf{rd}(D)$ for all concept descriptions C and D.

3. The role depth of a restriction is the successor of the role depth of the concept description in the restriction body, i.e. $\mathsf{rd}(Q\,r.\,C) := 1 + \mathsf{rd}(C)$ for all quantifiers $Q \in \{\exists, \forall, \geq n, \leq n\}$, role names $r \in N_R$ and concept descriptions C.
4. The role depth of a self restriction is just defined as 1, i.e. $\mathsf{rd}(\exists\, r.\,\mathsf{Self}) := 1$.

It is easy to see that the role-depth of a concept description is well-defined. However, equivalent concept descriptions do not neccessarily have the same role depth. For example the concept description $\bot$ and $\exists r.\bot$ are equivalent, but the former concept description has role depth 0 and the latter has role depth 1.

Definition 2 (Model-Based Most-Specific Concept Description). *Let (N_C, N_R) be a signature, $\mathcal{I} = (\Delta^{\mathcal{I}}, \cdot^{\mathcal{I}})$ an interpretation over (N_C, N_R), $\delta \in \mathbb{N}$ a role-depth bound, and $X \subseteq \Delta^{\mathcal{I}}$ a subset of the interpretation's domain. Then an $\mathcal{ALEQR}^{\mathsf{Self}}$-concept description C is called a* model-based most-specific concept description (mmsc) *of X w.r.t. $\mathcal{I}$ and δ, if it satisfies the following conditions.*

1. *C has a role depth of at most d, i.e. $\mathsf{rd}(C) \leq \delta$.*
2. *All elements of X are in the extension of C w.r.t. $\mathcal{I}$, i.e. $X \subseteq C^{\mathcal{I}}$.*
3. *For all concept descriptions D with $\mathsf{rd}(D) \leq \delta$ and $X \subseteq D^{\mathcal{I}}$ it holds that $C \sqsubseteq D$.*

Since all model-based most-specific concept descriptions of X w.r.t. $\mathcal{I}$ and δ are unique up to equivalence, we speak of *the* mmsc and denote it by $X^{\mathcal{I}_\delta}$.

Lemma 1. *Let $\mathcal{I}$ be an interpretation over the signature (N_C, N_R). Then the following statements hold for all subsets $X, Y \subseteq \Delta^{\mathcal{I}}$ and concept descriptions $C, D \in \mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R)$ with a role-depth $\leq \delta$.*

1. *$X \subseteq C^{\mathcal{I}}$ if and only if $X^{\mathcal{I}_\delta} \sqsubseteq C$.*
2. *$X \subseteq Y$ implies $X^{\mathcal{I}_\delta} \sqsupseteq Y^{\mathcal{I}_\delta}$.*
3. *$C \sqsubseteq D$ implies $C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$.*
4. *$X \subseteq X^{\mathcal{I}_\delta \mathcal{I}}$.*
5. *$C \sqsupseteq C^{\mathcal{I}\mathcal{I}_\delta}$.*
6. *$X^{\mathcal{I}_\delta} \equiv X^{\mathcal{I}_\delta \mathcal{I}\mathcal{I}_\delta}$.*
7. *$C^{\mathcal{I}} = C^{\mathcal{I}\mathcal{I}_\delta \mathcal{I}}$.*

It then follows that $\cdot^{\mathcal{I}\mathcal{I}_\delta}$ is a closure operator on the factorized concept poset $(\mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R), \sqsupseteq)$, and a concept inclusion $C \sqsubseteq D$ holds in $\mathcal{I}$, if and only if the implication $C \to D$ holds in the closure operator $\cdot^{\mathcal{I}\mathcal{I}_\delta}$. It follows that there is a canonical base of a concept inclusions holding in an interpretation $\mathcal{I}$, if $\mathcal{I}$ is finite.

Definition 3 (Least Common Subsumer). *Let C, D be $\mathcal{ALEQR}^{\mathsf{Self}}$-concept descriptions w.r.t. the signature (N_C, N_R). Then a concept description $E \in \mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R)$ is called a* least common subsumer *(lcs) of C and D, if the following conditions are fulfilled.*

1. *E subsumes both C and D, i.e. $C \sqsubseteq E$ and $D \sqsubseteq E$ hold.*
2. *Whenever F is a common subsumer of C and D, then F subsumes E, i.e. $C \sqsubseteq F$ and $D \sqsubseteq F$ implies $E \sqsubseteq F$ for all concept descriptions $F \in \mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R)$.*

It follows that least common subsumers are always unique up to equivalence. Hence we can speak of *the* lcs of two concept descriptions, and furthermore we denote it by $\mathsf{lcs}(C, D)$ or $C \sqcup D$. The definition can be canonically extended to an arbitrary number of concept descriptions, and then we write $\mathsf{lcs}(C_1, \ldots, C_n)$ or $\bigsqcup_{i=1}^{n} C_i$ for the least common subsumer of the concept descriptions $C_1, \ldots, C_n$.

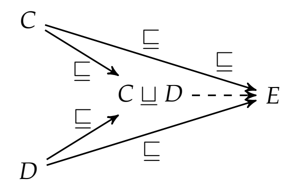

Fig. 3. The least common subsumer is a pullback in the category, whose objects are concept descriptions and whose morphisms are subsumptions.

Lemma 2. *Let $(X_t)_{t \in T}$ be a family of subsets $X_t \subseteq \Delta^{\mathcal{I}}$ and $(C_s)_{s \in S}$ a family of concept descriptions $C_s \in \mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R)$. Then the following statements hold.*

1. $(\bigcup_{t \in T} X_t)^{\mathcal{I}_\delta} \equiv \bigsqcup_{t \in T} X_t^{\mathcal{I}_\delta}$
2. $(\bigsqcap_{s \in S} C_s)^{\mathcal{I}} = \bigcap_{s \in S} C_s^{\mathcal{I}}$

Lemma 3. *If $C \sqsubseteq D$ holds in $\mathcal{I}$ and both C and D have a role depth $\leq \delta$, then also $C \sqsubseteq C^{\mathcal{I}\mathcal{I}_\delta}$ holds in $\mathcal{I}$, and $C \sqsubseteq D$ follows from $C \sqsubseteq C^{\mathcal{I}\mathcal{I}_\delta}$.*

Beforehand we have observed a pair of mappings that has similar properties like the well-known galois connection which is induced by a formal context. Consequently, we adapt the notions of formal concept and formal concept lattice as follows.

Definition 4 (Description Concept). *Let $\mathcal{I}$ be a finite interpretation over the signature (N_C, N_R) and $\delta \in \mathbb{N}$ a role-depth bound.*

A description concept *of $\mathcal{I}$ and δ is a pair (X, C), that consists of a subset $X \subseteq \Delta^{\mathcal{I}}$ and an $\mathcal{ALEQR}^{\mathsf{Self}}$-concept description over (N_C, N_R), such that X is the extension $C^{\mathcal{I}}$ and C is the model-based most-specific concept description $X^{\mathcal{I}_\delta}$. Furthermore we call X the* extent *and C the* intent *of (X, C). The set of all description concepts of $\mathcal{I}$ and δ is denoted as $\mathfrak{B}(\mathcal{I}, \delta)$. Analogously $\mathsf{Ext}(\mathcal{I}, \delta)$ and $\mathsf{Mmsc}(\mathcal{I}, \delta)$ denote the sets of all extents and intents, respectivelly.*

To ensure formal correctness, we require that $\mathfrak{B}(\mathcal{I}, \delta)$ only contains at most one description concept with the extent X. This is no limitation as we will see in the next lemma that all description concepts with the same extent have equivalent intents.

Definition 5 (Subconcept, Superconcept, Description Concept Lattice). *Let (X, C) and (Y, D) be two description concepts. Then (X, C) is a* subconcept *of (Y, D) if $X \subseteq Y$ holds. We then also write $(X, C) \leq (Y, D)$ and call (Y, D) a* superconcept *of (X, C).*

Additionally the pair $\underline{\mathfrak{B}}(\mathcal{I}, \delta) := (\mathfrak{B}(\mathcal{I}, \delta), \leq)$ is called description concept lattice *of $\mathcal{I}$ and δ.*

Lemma 4 (Order on Description Concepts). *Let $\mathcal{I}$ be a finite interpretation over the signature (N_C, N_R) and $\delta \in \mathbb{N}$ a role-depth bound.*

1. *For two description concepts (X, C) and (Y, D) it holds that*

$$(X, C) \leq (Y, D) \Leftrightarrow X \subseteq Y \Leftrightarrow C \sqsubseteq D.$$

2. *The relation $\leq$ is an order on $\mathfrak{B}(\mathcal{I}, \delta)$.*

We may furthermore observe that the set of all description concepts with the given order $\leq$ is a complete lattice.

Definition 6 (Description Lattice). *Let $\mathcal{I}$ be a finite interpretation over the signature (N_C, N_R) and $\delta \in \mathbb{N}$ a role-depth bound. Then $\underline{\mathfrak{B}}(\mathcal{I}, \delta)$ is a complete lattice, where the infima and suprema are given by the following equations.*

$$\bigwedge_{t \in T} (X_t, C_t) = \left(\bigcap_{t \in T} X_t, \left(\bigsqcap_{t \in T} C_t \right)^{\mathcal{I}\mathcal{I}_\delta} \right)$$

$$\bigvee_{t \in T} (X_t, C_t) = \left(\left(\bigcup_{t \in T} X_t \right)^{\mathcal{I}_\delta \mathcal{I}}, \underline{\bigsqcup_{t \in T}} C_t \right)$$

A description lattice is a nice visualization of the information provided in a description graph or in an interpretation, respectivelly. Since interpretations and descriptions are cryptomorphic definitions, we do not want to further distinguish between them. One can think of description lattices as a natural generalization of concept lattices, to not only allow conjunctions of attributes as intents, but also more complex concept descriptions, which are allowed in the underlying description logic. Of course, if the chosen description logic is $\mathcal{L}_0$, i.e. only allows for conjunctions $\sqcap$, then the concept lattices and description lattices w.r.t. $\mathcal{L}_0$ coincide. However, for more complex description logics like $\mathcal{EL}$ or $\mathcal{FLE}$ or extensions thereof, we can further involve roles in the intents of the description concepts, which adds further expressivity.

There is also a strong correspondence to the pattern structures and their lattices, as they have been introduced in [17]. The similarity operation is simply given by the least common subsumer mapping $\underline{\sqcup}$, which is the infimum in the lattice of all concept descriptions. Of course, the set of patterns consists of all concept descriptions that are expressible in the underlying description logic w.r.t. the given signature (N_C, N_R).

5 Induced Concept Contexts

Definition 7 (Induced Context). *Let $\mathcal{I}$ be an interpretation and $\mathcal{M}$ a set of concept descriptions, both over the signature (N_C, N_R). Then the* induced context *of $\mathcal{I}$ and $\mathcal{M}$ is defined as the formal context $\mathbb{K}_{\mathcal{I},\mathcal{M}} := (\Delta^{\mathcal{I}}, \mathcal{M}, I)$, where the incidence I is defined via $(x, C) \in I$ if and only if $x \in C^{\mathcal{I}}$. For a concept description C over (N_C, N_R) its* projection *to $\mathcal{M}$ is defined as $\pi_{\mathcal{M}}(C) := \{D \in \mathcal{M} \mid C \sqsubseteq D\}$. A concept description C is* expressible in terms of *$\mathcal{M}$, if $C \equiv \bigsqcap U$ holds for a subset $U \subseteq \mathcal{M}$. We have that $\bigsqcap \emptyset = \top$ and $\bigsqcap U = \bigsqcap_{C \in U} C$ hold for all subsets $\emptyset \neq U \subseteq \mathcal{M}$.*

Lemma 5. *Let $\mathcal{I}$ be an interpretation and $\mathcal{M}$ a set of concept descriptions. Then the following statements hold for all subsets $X, Y \subseteq \mathcal{M}$ and all concept descriptions C, D.*

1. $X \subseteq \pi_{\mathcal{M}}(C)$ *if and only if* $C \sqsubseteq \bigsqcap X$.
2. $X \subseteq Y$ *implies* $\bigsqcap X \sqsupseteq \bigsqcap Y$.
3. $C \sqsubseteq D$ *implies* $\pi_{\mathcal{M}}(C) \supseteq \pi_{\mathcal{M}}(D)$.
4. $X \subseteq \pi_{\mathcal{M}}(\bigsqcap X)$
5. $C \sqsubseteq \bigsqcap \pi_{\mathcal{M}}(C)$.
6. $\bigsqcap X \equiv \bigsqcap \pi_{\mathcal{M}}(\bigsqcap X)$.
7. $\pi_{\mathcal{M}}(C) = \pi_{\mathcal{M}}(\bigsqcap \pi_{\mathcal{M}}(C))$.

Lemma 6. *Let $\mathbb{K}_{\mathcal{I},\mathcal{M}}$ be an induced context. Then the following statements hold for all concept descriptions C over (N_C, N_R) and all subsets $U \subseteq \mathcal{M}$ and $X \subseteq \Delta^{\mathcal{I}}$.*

1. $\pi_{\mathcal{M}}(X^{\mathcal{I}}) = X^I$
2. $(\bigsqcap U)^{\mathcal{I}} = U^I$
3. $C^{\mathcal{I}} \subseteq \pi_{\mathcal{M}}(C)^I$
4. $\pi_{\mathcal{M}}\left((\bigsqcap U)^{\mathcal{I}\mathcal{I}}\right) = U^{II}$
5. $C \equiv \bigsqcap \pi_{\mathcal{M}}(C)$, *if C is expressible in terms of $\mathcal{M}$.*
6. $C^{\mathcal{I}} = \pi_{\mathcal{M}}(C)^I$, *if C is expressible in terms of $\mathcal{M}$.*
7. $U = \pi_{\mathcal{M}}(\bigsqcap U)$, *if U is an intent of $\mathbb{K}_{\mathcal{I},\mathcal{M}}$.*

The next lemma tells us that we can decide directly in the induced context $\mathbb{K}_{\mathcal{I},\mathcal{M}}$, whether a concept inclusion between conjunctions of concept descriptions of $\mathcal{M}$ holds in the given interpretation $\mathcal{I}$.

Lemma 7 (Implications and concept inclusions). *Let $\mathcal{I}$ be an interpretation and $\mathcal{M}$ a set of concept descriptions, both over the signature (N_C, N_R). Then for all subsets $X, Y \subseteq \mathcal{M}$, the concept inclusion $\bigsqcap X \sqsubseteq \bigsqcap Y$ holds in $\mathcal{I}$, if and only if the implication $X \to Y$ holds in $\mathbb{K}_{\mathcal{I},\mathcal{M}}$.*

Definition 8 (Approximation). *Let $\mathcal{I}$ be an interpretation over the signature (N_C, N_R), $\delta \in \mathbb{N}$ a role-depth bound, and $C \in \mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R)$ a concept description with its normal form $\bigsqcap_{A \in U} A \sqcap \bigsqcap_{(Q,r,D) \in \Pi} Q\, r.\, D$. Then the* approximation *of C w.r.t. $\mathcal{I}$ and δ is defined as the concept description*

$$\lfloor C \rfloor_{\mathcal{I},\delta} := \bigsqcap_{A \in U} A \sqcap \bigsqcap_{(Q,r,D) \in \Pi} Q\, r.\, D^{\mathcal{I}\mathcal{I}_\delta}.$$

Lemma 8. *For all concept descriptions C, D and role names r the following statements hold.*

1. $(C^{\mathcal{I}\mathcal{I}_\delta} \sqcap D)^{\mathcal{I}} = (C \sqcap D)^{\mathcal{I}}$.
2. $(Q\, r.\, C^{\mathcal{I}\mathcal{I}_\delta})^{\mathcal{I}} = (Q\, r.\, C)^{\mathcal{I}}$ *for all quantifiers $Q \in \{\exists, \forall, \geq n, \leq n\}$.*

Lemma 9. *For every interpretation $\mathcal{I}$ and every concept description C it holds that*

$$C^{\mathcal{I}\mathcal{I}_\delta} \sqsubseteq \lfloor C \rfloor_{\mathcal{I},\delta} \sqsubseteq C.$$

Lemma 10. *Let $\mathcal{I}$ be an interpretation and $\delta \in \mathbb{N}$ a role-depth bound, and define*

$$\mathcal{M}_{\mathcal{I},\delta} := \{\bot\} \cup \{A, \neg A \mid A \in N_C\} \cup \left\{ \begin{array}{l|l} \exists\, r.\, X^{\mathcal{I}_{\delta-1}}, & \\ \forall\, r.\, X^{\mathcal{I}_{\delta-1}}, & r \in N_R, \\ \geq n.\, r.\, X^{\mathcal{I}_{\delta-1}}, & 1 \leq m < n \leq \left|\Delta^{\mathcal{I}}\right|, \\ \leq m.\, r.\, X^{\mathcal{I}_{\delta-1}}, & \emptyset \neq X \subseteq \Delta^{\mathcal{I}} \\ \exists\, r.\, \mathsf{Self} & \end{array} \right\}.$$

Then every model-based most specific concept description of $\mathcal{I}$ with role-depth $\leq \delta$ is expressible in terms of $\mathcal{M}_{\mathcal{I},\delta}$. Furthermore the induced context *of $\mathcal{I}$ and δ is defined as the induced context $\mathbb{K}^{C}_{\mathcal{I},\delta} := \mathbb{K}_{\mathcal{I},\mathcal{M}_{\mathcal{I},\delta}}$ of $\mathcal{I}$ and $\mathcal{M}_{\mathcal{I},\delta}$.*

Lemma 11 (Intents and MMSCs). *Let $\mathcal{I}$ be an interpretation over (N_C, N_R) and $\mathbb{K}_{\mathcal{I},\delta}$ its induced context w.r.t. the role-depth bound $\delta \in \mathbb{N}$. Then the following statements hold for all subsets $U \subseteq \mathcal{M}_{\mathcal{I},\delta}$ and concept descriptions C over (N_C, N_R).*

1. $(\bigsqcap U)^{\mathcal{I}\mathcal{I}_\delta} \equiv \bigsqcap U^{II}$.
2. *If U is an intent of $\mathbb{K}_{\mathcal{I},\delta}$, then $\bigsqcap U$ is a mmsc of $\mathcal{I}$ with role-depth $\leq \delta$.*
3. *If C is a mmsc of $\mathcal{I}$ with role-depth $\leq \delta$, then $\pi_{\mathcal{M}_{\mathcal{I},\delta}}(C)$ is an intent of $\mathbb{K}_{\mathcal{I},\delta}$.*

As a consequence, the mapping $\bigsqcap\colon \mathcal{M}_{\mathcal{I},\delta} \to \mathcal{ALEQR}^{\mathsf{Self}}(N_C, N_R)$ is an isomorphism from the intent-lattice $(\mathsf{Int}(\mathbb{K}_{\mathcal{I},\delta}), \cap)$ to the mmsc-lattice $(\mathsf{Mmsc}(\mathcal{I},\delta), \underline{\sqcup})$, and has the inverse $\pi_{\mathcal{M}_{\mathcal{I},\delta}}$. This shows the strong correspondence between the formal concept lattice of $\mathbb{K}_{\mathcal{I},\delta}$ and the description concept lattice of $\mathcal{I}$w.r.t. role depth $\leq \delta$. We can infer the following corollary from lemmata 11 and 5.

Corollary 1. *The intent lattice of $\mathbb{K}_{\mathcal{I},\delta}$ is isomorphic to the mmsc lattice of $\mathcal{I}, \delta$.*

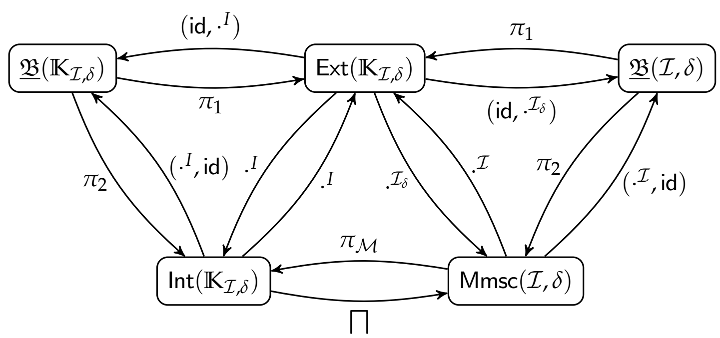

Fig. 4. Overview on the isomorphisms between the extent lattice, intent lattice and mmsc lattice of $\mathbb{K}_{\mathcal{I},\delta}$ and $\mathcal{I},\delta$, respectivelly. Note that $\mathsf{Ext}(\mathbb{K}_{\mathcal{I},\delta}) = \mathsf{Ext}(\mathcal{I},\delta)$ holds.

We can further observe that the concept inclusions holding in $\mathcal{I}$ and the implications holding in $\mathbb{K}_{\mathcal{I},\delta}$ are also in a strong correspondence. We can show that whenever the implication $U \to V$ holds in $\mathbb{K}_{\mathcal{I},\delta}$, then also the concept inclusion $\bigsqcap U \sqsubseteq \bigsqcap V$ holds in $\mathcal{I}$. Furthermore, since every mmsc of $\mathcal{I}$ with a role depth $\leq \delta$ is expressible in terms of $\mathcal{M}_{\mathcal{I},\delta}$, and conjunctions of intents of $\mathbb{K}_{\mathcal{I},\delta}$ are exactly the mmscs of $\mathcal{I}$, and every concept inclusion $C \sqsubseteq D$ holding in $\mathcal{I}$ is entailed by the concept inclusion $C \sqsubseteq C^{\mathcal{I}\mathcal{I}}$, we can deduce that indeed every concept inclusion holding in $\mathcal{I}$ is entailed by the transformation of the canonical implicational base of $\mathbb{K}_{\mathcal{I},\delta}$, which consists of all CGIs that have a conjunction of a pseudo-intent as premise and the conjunction of the closure of the pseudo-intent as conclusion.

Lemma 12. *Let $\mathcal{I}$ be an interpretation over the signature (N_C, N_R), $\delta \in \mathbb{N}$ a role-depth bound and $C \sqsubseteq D$ be a concept inclusion, such that both concepts C, D have a role-depth $\leq \delta$.*

1. *If D is expressible in terms of $\mathcal{M}_{\mathcal{I},\delta}$ and the implication $\pi_{\mathcal{M}_{\mathcal{I},\delta}}(C) \to \pi_{\mathcal{M}_{\mathcal{I},\delta}}(D)$ holds in $\mathbb{K}_{\mathcal{I},\delta}$, then the concept inclusion $C \sqsubseteq D$ holds in $\mathcal{I}$.*
2. *If C is expressible in terms of $\mathcal{M}_{\mathcal{I},\delta}$ and the concept inclusion $C \sqsubseteq D$ holds in $\mathcal{I}$, then the implication $\pi_{\mathcal{M}_{\mathcal{I},\delta}}(C) \to \pi_{\mathcal{M}_{\mathcal{I},\delta}}(D)$ holds in $\mathbb{K}_{\mathcal{I},\delta}$.*

Corollary 2 (Concept Inclusion Base). *Let $\mathcal{I}$ be an interpretation over the signature (N_C, N_R) and $\delta \in \mathbb{N}$ a role-depth bound. Then the following statements hold:*

1. *For all subsets $X, Y \subseteq \mathcal{M}_{\mathcal{I},\delta}$, the implication $X \to Y$ holds in $\mathbb{K}_{\mathcal{I},\delta}$, if and only if the concept inclusion $\bigsqcap X \sqsubseteq \bigsqcap Y$ holds in $\mathcal{I}$.*
2. *The intents of $\mathbb{K}_{\mathcal{I},\delta}$ are exactly the model-based most-specific concept descriptions of $\mathcal{I}$ with role-depth bound $\leq \delta$.*
3. *If $\mathcal{L}$ is an implicational base for $\mathbb{K}_{\mathcal{I},\delta}$, then $\bigsqcap \mathcal{L} := \{\bigsqcap X \sqsubseteq \bigsqcap Y \mid X \to Y \in \mathcal{L}\}$ is a sound and complete TBox for all concept inclusions holding in $\mathcal{I}, \delta$. Especially this holds for the following TBox.*

$$\left\{ \bigsqcap P \sqsubseteq \bigsqcap P^{II} \;\middle|\; P \text{ is a pseudo-intent of } \mathbb{K}_{\mathcal{I},\delta} \right\}$$

6 Induced Role Contexts

Role contexts have been introduced by Zickwolff [31] and have been used by Rudolph [27] for gaining knowledge on binary relations or roles (that are interpreted as binary relations). We use their definition here for the deduction of complex role inclusions holding in an interpretation.

Definition 9 (Induced Role Context). *Let $\mathcal{I}$ be an interpretation over the signature (N_C, N_R) and $\delta \in \mathbb{N}$ a role depth bound. Furthermore assume, that $X = \{x_0, x_1, \ldots, x_\delta\}$ is a set of $\delta + 1$* variables. *Then the* induced role context *for $\mathcal{I}$ and δ is defined as*

$$\mathbb{K}^R_{\mathcal{I},\delta} := \left(\left(\Delta^{\mathcal{I}}\right)^X, X \times N_R \times X, J \right),$$

where $(f, (x, r, y)) \in J$ if and only if $(f(x), f(y)) \in r^{\mathcal{I}}$ holds.

Lemma 13 (Role Inclusions and Implications). *Let $\mathcal{I}$ be an interpretation over (N_C, N_R), $\delta \in \mathbb{N}$ a role-depth bound and $n \leq \delta$. Then the complex role inclusion $r_1 \circ r_2 \circ \ldots \circ r_n \sqsubseteq s$ holds in $\mathcal{I}$, if and only if the implication $\{(x_0, r_1, x_1), (x_1, r_2, x_2), \ldots, (x_{n-1}, r_n, x_n)\} \to \{(x_0, s, x_n)\}$ holds in the induced role context $\mathbb{K}^R_{\mathcal{I},\delta}$.*

We define a constraining closure operator ϕ_R on the attribute set $X \times N_R \times X$ of the induced role context as follows. Since we are only interested in implications, whose premise contains a subset of the form $\{(x_0, r_1, x_1), (x_1, r_2, x_2), \ldots, (x_{k-1}, r_k, x_k)\}$ we call all subsets that contain such a set closed. Formally we thus define

$$\phi_R(B) := \begin{cases} B & \begin{array}{l} \text{if } \exists k \in \mathbb{N}_+ \exists r_1, r_2, \ldots, r_k \in N_R \\ \quad \exists \{x_0, x_1, \ldots, x_k\} \in \binom{X}{k+1} : \\ \quad \{(x_0, r_1, x_1), \ldots, (x_{k-1}, r_k, x_k)\} \subseteq B, \end{array} \\ B \cup \{(x_0, r, x_1) \mid r \in N_R\} & \text{otherwise.} \end{cases}$$

Then ϕ_R is extensive, monotone and idempotent, i.e. a closure operator.

Theorem 1 (Role Inclusion Base). *Let $\mathcal{I}$ be an interpretation over (N_C, N_R). If $\mathcal{L}$ is a ϕ_R-constrained implicational base of $\mathbb{K}_{\mathcal{I},R}$, then the following RBox $\mathcal{R}_{\mathcal{I},\delta}$ is sound, complete and irredundant for all complex role inclusions holding in $\mathcal{I}, \delta$.*

$$\mathcal{R}_{\mathcal{I},\delta} := \left\{ r_1 \circ r_2 \circ \ldots \circ r_k \sqsubseteq s \;\middle|\; \begin{array}{l} \exists X \to Y \in \mathcal{L} \\ \exists r_1, r_2, \ldots, r_k, s \in N_R \\ \exists \{x_1, x_2, \ldots, x_{k+1}\} \in \binom{X}{k+1} : \\ X \supseteq \{(r_1, x_1, x_2), (r_2, x_2, x_3), \ldots, (r_k, x_k, x_{k+1})\} \\ Y \ni (s, x_1, x_{k+1}) \end{array} \right\}$$

7 Construction of the Knowledge Base

By means of the results of the previous sections 5 and 6 we are now ready to formulate a knowledge base for an interpretation $\mathcal{I}$ or for a description graph $\mathcal{G}$, respectivelly. Beforehand, it is neccessary to inspect the interplay of role and concept inclusions to ensure irredundancy of the knowledge base. At first we list some trivial concept inclusions that hold in all interpretations.

Lemma 14. *Let $m, n \in \mathbb{N}_+$ be non-negative integers with $n < m$, $r \in N_R$ a role name and C a concept description. The following general concept inclusions hold in every interpretation $\mathcal{I}$.*

$$\begin{aligned} A \sqcap \neg A &\sqsubseteq \bot \\ \exists r.\,\mathsf{Self} \sqcap \forall r.\,C &\sqsubseteq C \\ \exists r.\,\mathsf{Self} \sqcap C &\sqsubseteq \exists r.\,C \\ \exists r.\,\mathsf{Self} \sqcap C \sqcap \leq 1.\,r.\,C &\sqsubseteq \forall r.\,C \\ \exists r.\,C \sqcap \forall r.\,D &\sqsubseteq \exists r.\,(C \sqcap D) \\ \geq n.\,r.\,C \sqcap \forall r.\,D &\sqsubseteq \geq n.\,r.\,(C \sqcap D) \\ \leq n.\,r.\,C \sqcap \forall r.\,D &\sqsubseteq \leq n.\,r.\,(C \sqcap D) \\ \exists r.\,C &\sqsubseteq \geq 1.\,r.\,C \\ \geq n.\,r.\,C &\sqsubseteq \exists r.\,C \\ \leq n.\,r.\,C &\sqsubseteq \leq m.\,r.\,C \\ \geq m.\,r.\,C &\sqsubseteq \geq n.\,r.\,C \\ \geq \left|\Delta^{\mathcal{I}}\right|.\,r.\,C &\sqsubseteq C \sqcap \forall r.\,C \sqcap \exists r.\,\mathsf{Self} \\ \top &\sqsubseteq \leq \left|\Delta^{\mathcal{I}}\right|.\,r.\,C \end{aligned}$$

Please note that there are no direct subsumptions between existential restrictions $\exists r.\,C$ and value restrictions $\forall r.\,C$, i.e. both $\exists r.\,C \sqsubseteq \forall r.\,C$ and $\forall r.\,C \sqsubseteq \exists r.\,C$ do not hold. There is also a crossover between both constructors existential restriction and value restriction. The constructor is denoted by $\forall\!\!\exists$ and has the semantics $(\forall\!\!\exists\, r.\,C)^{\mathcal{I}} := (\exists r.\,C)^{\mathcal{I}} \cap (\forall r.\,C)^{\mathcal{I}}$, i.e. a domain element is in the extension of $\forall\!\!\exists\, r.\,C$, iff there is an r-successor in C and all r-successors are in C.

The next two lemmata show us, which concept inclusions can be inferred from known role inclusions.

Lemma 15. *Let $\mathcal{I}$ be a model of the role inclusion axiom $r \sqsubseteq s$, and C an arbitrary concept description. Then $\mathcal{I}$ is also a model of the following general concept inclusions.*

$$\begin{aligned} Q_1\, r.\, C &\sqsubseteq Q_1\, s.\, C \\ \exists\, r.\, \mathsf{Self} &\sqsubseteq \exists\, s.\, \mathsf{Self} \\ Q_2\, s.\, C &\sqsubseteq Q_2\, r.\, C \end{aligned}$$

The quantifiers can be arbitrarily chosen as $Q_1 \in \{\exists, \geq n\}$ and $Q_2 \in \{\forall, \leq n\}$, respectivelly, and n is a non-negative integer from $\mathbb{N}_+$.

Lemma 16. *Let $\mathcal{I}$ be a model of the complex role inclusion $r_1 \circ r_2 \circ \ldots \circ r_k \sqsubseteq s$, and C an arbitrary concept description. Then $\mathcal{I}$ is also a model of the following concept inclusions.*

$$\begin{aligned} \exists\, r_1.\, \exists\, r_2. \ldots Q_1\, r_k.\, C &\sqsubseteq Q_1\, s.\, C \\ Q_2\, s.\, C &\sqsubseteq \forall\, r_1.\, \forall\, r_2. \ldots Q_2\, r_k.\, C \end{aligned}$$

The quantifiers can be arbitrarily chosen as $Q_1 \in \{\exists, \geq n\}$ and $Q_2 \in \{\forall, \leq n\}$, respectivelly, and n is a non-negative integer from $\mathbb{N}_+$.

As final step we use the trivial concept inclusions and concept inclusions that are entailed by valid role inclusions to define some background knowledge for the computation of the canonical implicational base of the induced concept context, which is trivial in terms of description logics, but not for formal concept analysis, due to their different semantics.

Theorem 2 (Knowledge Base). *Let $\mathcal{I}$ be an interpretation over the signature (N_C, N_R) and $\delta \in \mathbb{N}$ a role-depth bound. Furthermore assume that $\mathcal{L}$ is an implicational base of the induced concept context $\mathbb{K}^C_{\mathcal{I},\delta}$ w.r.t. the background knowledge*

$$\begin{aligned} \mathcal{S}_{\mathcal{I}} := &\left\{ \{C\} \to \{D\} \,\middle|\, \begin{array}{l} C, D \in \mathcal{M}_{\mathcal{I},\delta}, \\ C \sqsubseteq D \end{array} \right\} \\ &\cup \{\{A, \neg A\} \to \mathcal{M}_{\mathcal{I},\delta} \mid A \in N_C\} \\ &\cup \left\{ \begin{array}{r} \left\{\exists\, r.\, X^{\mathcal{I}_{\delta-1}}, \forall\, r.\, Y^{\mathcal{I}_{\delta-1}}\right\} \to \left\{\exists\, r.\, Z^{\mathcal{I}_{\delta-1}}\right\}, \\ \left\{\geq n.\, r.\, X^{\mathcal{I}_{\delta-1}}, \forall\, r.\, Y^{\mathcal{I}_{\delta-1}}\right\} \to \left\{\geq n.\, r.\, Z^{\mathcal{I}_{\delta-1}}\right\}, \\ \left\{\leq m.\, r.\, X^{\mathcal{I}_{\delta-1}}, \forall\, r.\, Y^{\mathcal{I}_{\delta-1}}\right\} \to \left\{\leq m.\, r.\, Z^{\mathcal{I}_{\delta-1}}\right\}, \end{array} \,\middle|\, \begin{array}{l} r \in N_R, \\ \emptyset \neq X, Y, Z \subseteq \Delta^{\mathcal{I}}, \\ Z^{\mathcal{I}_{\delta-1}} \equiv X^{\mathcal{I}_{\delta-1}} \sqcap Y^{\mathcal{I}_{\delta-1}}, \\ 1 \leq m < n \leq \left|\Delta^{\mathcal{I}}\right| \end{array} \right\} \\ &\cup \left\{ \begin{array}{c} \left\{\exists\, r.\, X^{\mathcal{I}_{\delta-1}}\right\} \to \left\{\exists\, s.\, X^{\mathcal{I}_{\delta-1}}\right\}, \\ \left\{\forall\, s.\, X^{\mathcal{I}_{\delta-1}}\right\} \to \left\{\forall\, r.\, X^{\mathcal{I}_{\delta-1}}\right\}, \\ \left\{\geq n.\, r.\, X^{\mathcal{I}_{\delta-1}}\right\} \to \left\{\geq n.\, s.\, X^{\mathcal{I}_{\delta-1}}\right\}, \\ \left\{\leq m.\, s.\, X^{\mathcal{I}_{\delta-1}}\right\} \to \left\{\leq m.\, r.\, X^{\mathcal{I}_{\delta-1}}\right\}, \\ \{\exists\, r.\, \mathsf{Self}\} \to \{\exists\, s.\, \mathsf{Self}\} \end{array} \,\middle|\, \begin{array}{l} r, s \in N_R, \\ r \sqsubseteq s \in \mathcal{R}, \\ 1 \leq m < n \leq \left|\Delta^{\mathcal{I}}\right|, \\ \emptyset \neq X \subseteq \Delta^{\mathcal{I}} \end{array} \right\}. \end{aligned}$$

Then $\mathcal{K}_{\mathcal{I},\delta} = (\mathcal{T}_{\mathcal{I},\delta}, \mathcal{R}_{\mathcal{I},\delta})$ is a knowledge base for $\mathcal{I}$, where $\mathcal{T}_{\mathcal{I},\delta} := \bigsqcap \mathcal{L}$ holds as in corollary 2 and $\mathcal{R}_{\mathcal{I},\delta}$ is defined as in theorem 1.

8 Other Description Logics

If only a lower expressivity of the underlying description logic is neccessary, then one could also use $\mathcal{EL}$, $\mathcal{FLE}$ or extensions thereof with role hierarchies $\mathcal{H}$ or complex role inclusions $\mathcal{R}$. All of the previous results are still valid, however one has to remove some of the used concept descriptions that are not expressible in the chosen description logic. Figure 5 gives an overview on description logics that have a lower expressivity than $\mathcal{ALEQR}^{\mathsf{Self}}$ and could also be used for knowledge acquisition.

constructor	$\mathcal{EL}$	$\mathcal{FL}_0$	$\mathcal{FLE}$	$\mathcal{ALE}$	$\mathcal{Q}$	Self	$\mathcal{H}$	$\mathcal{R}$
$\bot$				×				
$\top$	×	×	×	×				
$\neg A$				×				
$C \sqcap D$	×	×	×	×				
$\exists r.C$	×		×	×	×			
$\forall r.C$		×	×	×				
$\geq n.r.C$					×			
$\leq n.r.C$					×			
$\exists r.\mathsf{Self}$						×		
$C \sqsubseteq D$	×	×	×	×				
$C \equiv D$	×	×	×	×				
$r \sqsubseteq s$							×	×
$r_1 \circ \ldots \circ r_n \sqsubseteq s$								×

Fig. 5. Overview on various Description Logics below $\mathcal{ALEQR}^{\mathsf{Self}}$

8.1 Role Hierarchies $\mathcal{H}$ instead of Complex Role Inclusions $\mathcal{R}$

In the special case of simple role inclusions provided by the extension $\mathcal{H}$ it is not neccessary to use the induced role context. We can directly extract the role hierarchy from the interpretation $\mathcal{I}$ or the description graph $\mathcal{G}$, respectivelly, as follows.

First, we want to extract a minimal RBox $\mathcal{R}_{\mathcal{I}}$ from the interpretation that entails all role inclusion axioms holding in $\mathcal{I}$. We therefore define an equivalence relation $\equiv_{\mathcal{I}}$ on the role names as follows: $r \equiv_{\mathcal{I}} s$ if and only if $r^{\mathcal{I}} = s^{\mathcal{I}}$. Then let $N_R^{\mathcal{I}}$ be a set of representatives of this equivalence relation, i.e. $\left| N_R^{\mathcal{I}} \cap [r]_{\equiv_{\mathcal{I}}} \right| = 1$ for all role names $r \in N_R$. Then add the following role equivalence axioms to $\mathcal{R}_{\mathcal{I}}$: For each representant role $r \in N_R^{\mathcal{I}}$ add the axioms $r \equiv s$ for all $s \in [r]_{\equiv_{\mathcal{I}}} \setminus \{r\}$. Then furthermore define an order relation $\sqsubseteq_{\mathcal{I}}$ on the representants $N_R^{\mathcal{I}}$ by $r \sqsubseteq_{\mathcal{I}} s$ iff $r^{\mathcal{I}} \subseteq s^{\mathcal{I}}$. Let $\prec_{\mathcal{I}}$ be the neighborhood relation of $\sqsubseteq_{\mathcal{I}}$, then add the role inclusion axioms $r \sqsubseteq s$ for each pair $r \prec_{\mathcal{I}} s$ to the RBox $\mathcal{R}_{\mathcal{I}}$. Then the constructed RBox is obviously minimal w.r.t. the property to entail all valid role inclusion axioms holding in the interpretation $\mathcal{I}$. Then the RBox in $\mathcal{K}_\mathcal{I}$ is simply defined as follows.

$$\mathcal{R}_\mathcal{I} := \left\{ r \equiv s \;\middle|\; r \in N_R^\mathcal{I}, s \in [r]_{\equiv_\mathcal{I}} \setminus \{r\} \right\} \cup \left\{ r \sqsubseteq s \;\middle|\; r, s \in N_R^\mathcal{I}, r \prec_\mathcal{I} s \right\}$$

9 Conclusion

We have provided an extension of the results of Distel [13] for the deduction of knowledge bases from interpretations in the more expressive description logic $\mathcal{ALEQR}^{\mathsf{Self}}$. Since role-depth-bounded model-based most-specific concept descriptions always exist, this technique can always be applied. Furthermore the computation of the knowledge base has been reduced to the computation of implicational bases of formal contexts, which is a well-understood problem that has several available algorithms. One could for example use the standard NextClosure algorithm by Ganter [16], or the parallel algorithm for the computation of the canonical base that has been introduced in [24] and implemented in [21]. The presented methods are prototypically implemented in Concept Explorer FX [21].

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