

2. Exercises for the Course „Description Logics“

Exercise 6:

Let $G = (V, E)$ be a directed graph represented as a set of PROLOG facts

$$\{\text{directly-connected}(x, y) \mid (x, y) \in E\}.$$

Consider three PROLOG programs that compute whether two nodes of a graph are connected:

- (a) $\text{connected}(x, y) :- \text{directly-connected}(x, y)$
 $\text{connected}(x, y) :- \text{directly-connected}(x, z), \text{connected}(z, y)$
- (b) $\text{connected}(x, y) :- \text{directly-connected}(x, y)$
 $\text{connected}(x, y) :- \text{connected}(z, y), \text{directly-connected}(x, z)$
- (c) $\text{connected}(x, y) :- \text{directly-connected}(x, z), \text{connected}(z, y)$
 $\text{connected}(x, y) :- \text{directly-connected}(x, y)$

Do the following:

- For each of the three programs, determine whether it is sound, complete, and terminating.
- Rewrite each program as a set of implication in first-order logic. Are the three sets logically equivalent?
- A KR formalism is *declarative* if the meaning of its terms is defined independently of a concrete interpreter or reasoning algorithm. Is KR in PROLOG declarative?

Exercise 7:

Let α and β be propositional formulae. Prove or disprove the following propositions:

- (a) If $\varphi \rightarrow \psi$ and φ are valid, then ψ is valid.
- (b) If $\varphi \rightarrow \psi$ and φ are satisfiable, then ψ is satisfiable.
- (c) If $\varphi \rightarrow \psi$ is valid and φ is satisfiable, then ψ is satisfiable.

Exercise 8:

A propositional formula using only the constructors \wedge , \vee , and \neg is in *negation normal form* (NNF) if negation occurs only in front of propositional variables.

Prove that each propositional formula can be transformed into an equivalent one in NNF.

Exercise 9:

Define a generic frame that describes the prototypical object “computer science course”. Use slots

- *Title*,
- *Lecturer*,
- *Type of course*, and
- *Hours per week*.

Find other meaningful slots. Then construct an instance frame for the generic frame.