

Selected Topics in Automata and Logic

Exercise Sheet 11

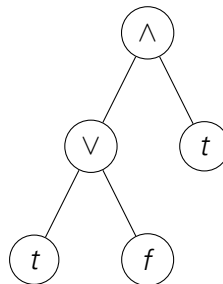
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Summer Semester 2010

Exercise 1

- Show that all finite tree languages are regular.
- Show that all finite tree languages are recognizable by a tree walking automaton.

Exercise 2

Consider boolean expressions over the binary operators \wedge and \vee and the nullary symbols t and f . These expressions can be represented by trees, e. g. $b = (t \vee f) \wedge t$ can be represented by



Construct tree walking automata that accept the following tree languages over the alphabet $\Sigma = \{\wedge, \vee, t, f\}$:

- the language of all trees that have at least one leaf that is labelled with f , and
- the language of all trees that represent well formed boolean expressions, i. e. where all leafs are labelled with t or f and all internal nodes are labelled with \vee or \wedge , and
- the language of all trees that represent a boolean expression that evaluates to true.

For each of your automata give an accepting run for the tree representing b .

Exercise 3

Give FOL+TC formulae that describe the languages from Exercise 2.

Exercise 4

Let Σ be an alphabet. For each tree t over Σ we define a language $\text{path}(t)$ as follows.

$$\text{path}(t) := \{a_1 a_2 \dots a_n \in \Sigma^* \mid \text{there is a leaf } i_1 i_2 \dots i_n \text{ in } t \\ \text{such that } t(i_1 \dots i_k) = a_k \text{ for all } k \leq n\},$$

i. e. $\text{path}(t)$ is the language of all words labelling a path in t . For a tree language B define

$$\text{path}(B) := \bigcup_{t \in B} \text{path}(t).$$

Conversely, for a language $L \subseteq \Sigma^*$ let

$$\text{tree}(L) := \{t \in T_\Sigma \mid \text{path}(t) \subseteq L\}.$$

Prove or refute.

- $\text{tree}(\text{path}(B)) = B$
- $\text{path}(\text{tree}(L)) = L$
- If L is regular then $\overline{\text{tree}(L)}$ is recognizable by a 1-head tree walking automaton.
- If $\text{path}(L)$ is regular then L is recognizable by a 1-head tree walking automaton.

Exercise 5

Let $\mathcal{L}_{1\text{TWA}}$ be the set of all tree languages over a fixed alphabet Σ that can be recognized by a 1-head tree walking automaton. Prove or disprove.

- $\mathcal{L}_{1\text{TWA}}$ is closed under union.
- $\mathcal{L}_{1\text{TWA}}$ is closed under intersection.