

Faculty of Computer Science Institute of Theoretical Computer Science, Chair of Automata Theory

# **Description Logics**

#### **Exercise Sheet 4**

Dr. Anni-Yasmin Turhan / Dipl.-Math. Felix Distel Summer Semester 2011

#### **Exercise 1**

In the lecture we saw that bisimulations allow to compare the expressive powers of DLs.

- a) Extend the notion of bisimulation relation to  $\mathcal{ALCN}$ .
- b) Prove that ALCN is bisimulation invariant for the bisimulation relation defined in (a).
- c) Prove that  $\mathcal{ALCQ}$  is more expressive than  $\mathcal{ALCN}$ .

## Exercise 2

Let  $\mathcal{K} = (\mathcal{T}, \mathcal{A})$  be a consistent knowledge base. We write  $C \sqsubseteq_{\mathcal{K}} D$  if  $C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$  for all models  $\mathcal{I}$  of  $\mathcal{K}$ . Prove that for all  $\mathcal{ALC}$ -concepts C and D, we have  $C \sqsubseteq_{\mathcal{K}} D$  iff  $C \sqsubseteq_{\mathcal{T}} D$ . Hint: Use disjoint unions.

#### Exercise 3

Show the following claim: If a concept *C* is satisfiable w.r.t. an  $\mathcal{ALC}$ -TBox  $\mathcal{T}$ , then for all  $n \ge 1$  there is a model  $\mathcal{I}_n$  of  $\mathcal{T}$  such that:  $|\mathcal{C}^{\mathcal{I}_n}| \ge n$ .

## Exercise 4

Prove or refute the following claim:

Given an  $\mathcal{ALC}$ -concept C and an  $\mathcal{ALC}$ -TBox  $\mathcal{T}$ . If  $\mathcal{I}$  is an interpretation and  $\mathcal{J}$  its filtration w.r.t. Sub(C)  $\cup$  Sub( $\mathcal{T}$ ), then the relation  $\varrho = \{(d, [d]) \mid d \in \Delta^{\mathcal{I}}, [d] \in \Delta^{\mathcal{J}}, d \simeq [d]\}$  is a bisimulation.