



Description Logics

Exercise Sheet 14

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Exercise 1

Prove that for a given TBox \mathcal{T} and a consequence $A \sqsubseteq B$ there can be exponentially many MinAs for $A \sqsubseteq B$. Hint: You do not need any concept constructors except conjunction.

Exercise 2

Complete the proof of Theorem 8.5 for the case where ψ is obtained using (R3).

Exercise 3

Let \mathcal{L}_1 and \mathcal{L}_2 be two DL-languages. We define the \mathcal{L}_2 -lcs as follows. Let C_1, C_2 be two \mathcal{L}_1 -concept descriptions. An \mathcal{L}_2 -concept description C is called the \mathcal{L}_2 -lcs of C_1 and C_2 iff

- (1) $C_1 \sqsubseteq C$ and $C_2 \sqsubseteq C$, and
- (2) for all \mathcal{L}_2 -concept descriptions D it holds that $C_1 \sqsubseteq D$ and $C_2 \sqsubseteq D$ imply $C \sqsubseteq D$.

This generalizes the standard notion of least common subsumers since it allows different logics for C_1, C_2 and their lcs.

Consider the following $\mathcal{FL}\mathcal{E}$ -concept descriptions

$$C_1 = \exists r.C \sqcap \exists r.D \sqcap \forall r.(A \sqcap B),$$
$$C_2 = \exists r.B \sqcap \exists r.D \sqcap \forall r.(A \sqcap D).$$

Find

- a) the \mathcal{EL} -lcs,
- b) the $\mathcal{FL}\mathcal{E}$ -lcs¹, and
- c) the \mathcal{ALC} -lcs

of C_1 and C_2 . Hint: It is not necessary to provide an algorithm that can compute the lcs in the general case.

¹ $\mathcal{FL}\mathcal{E}$ provides conjunction, existential restrictions and value restrictions.