

Faculty of Computer Science Institute of Theoretical Computer Science, Chair of Automata Theory

Fuzzy Logic

Solutions to Exercise Sheet 4

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Exercise 1

- a)
 $$\begin{split} \Pi &\vdash ((\varphi \to \neg \varphi) \& \varphi) \to \mathbf{0} \quad (\Pi 2) + (A3) \\ \Pi &\vdash (\varphi \to \neg \varphi) \to (\varphi \to \mathbf{0}) \quad (A6) \\ \Pi &\vdash (\varphi \to (\varphi \to \mathbf{0})) \to \neg \varphi \quad \text{def. of } \neg \\ \Pi &\vdash ((\varphi \& \varphi) \to \mathbf{0}) \to \neg \varphi \quad (A6) + (A1) \\ \Pi &\vdash \neg (\varphi \& \varphi) \to \neg \varphi \quad \text{def. of } \neg \end{split}$$
- c) For the first direction it only remains to show that $\Pi \vdash \neg \neg \neg \varphi \rightarrow \neg \varphi$. From the properties of negation in BL we obtain
 - $$\begin{split} &\Pi\vdash\varphi\to\neg\neg\varphi\text{, and}\\ &\Pi\vdash(\varphi\to\neg\neg\varphi)\to(\neg\neg\neg\varphi\to\neg\varphi)\text{, and thus}\\ &\Pi\vdash\neg\neg\neg\varphi\to\neg\varphi. \end{split}$$

Consider the other direction where we want to prove $\Pi'' \vdash (\varphi \land \neg \varphi) \to \mathbf{0}$. $\Pi''' \vdash \neg \neg \varphi \to ((\varphi \& \varphi \to \underbrace{\varphi \& \mathbf{0}}_{=\mathbf{0}}) \to \underbrace{(\varphi \to \mathbf{0})}_{\neg \varphi}), \quad (\Pi 1)$ $\Pi''' \vdash \neg \varphi \to ((\varphi \& \varphi \to \mathbf{0}) \to \neg \varphi), \quad (L. 2.14 (1))$

 $\begin{array}{l} \Pi''' \vdash \neg \varphi \rightarrow ((\varphi \& \varphi \rightarrow \mathbf{0}) \rightarrow \neg \varphi), \quad (L. 2.14 (1)) \\ \Pi''' \vdash (\neg \varphi \lor \neg \neg \varphi) \rightarrow ((\varphi \& \varphi \rightarrow \mathbf{0}) \rightarrow \neg \varphi), \quad (\text{properties of weak disjunction}) \\ \Pi''' \vdash ((\varphi \& \varphi \rightarrow \mathbf{0}) \rightarrow \neg \varphi), \quad \text{c}) + \text{mod. pon.} \\ \text{The rest follows from a).} \end{array}$