



## Automata and Logic

### Exercise Sheet 8

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#### Exercise 37

Let  $\Sigma = \{a, b\}$ , and  $L \subseteq \Sigma^*$  be defined by the regular expression  $(a^*bb^*)^*$ . Show that

$$\lim L = \{\alpha \in \Sigma^\omega \mid \text{if } \alpha(i, i) = a, \text{ then there is a } j > i \text{ with } \alpha(j, j) = b\}.$$

#### Exercise 38

Give Büchi automata that recognise the following  $\omega$ -regular languages over the alphabet  $\Sigma := \{a, b, c\}$ :

- $L_1 := \{\alpha \in \Sigma^\omega \mid \exists i \in \mathbb{N}: \alpha(i, i+2) = abc\}$ ;
- $L_2 := \{\alpha \in \Sigma^\omega \mid \{i \in \mathbb{N} \mid \alpha(i, i+2) = abc\} \text{ is infinite}\}$ ; and
- $L_3 := (a^+b^+c^+)^\omega$ .

#### Exercise 39

- Show that the construction used in the proof of Lemma 4.7.1 does not work for automata whose initial state is reachable from another state.
- Complete the proof of Lemma 4.7 from the lecture by showing the following:

If  $L_1, L_2 \subseteq \Sigma^\omega$  are Büchi recognisable, then  $L_1 \cup L_2$  is Büchi recognisable.

#### Exercise 40

Let  $\Sigma$  be an alphabet, and  $L, L_1, L_2 \subseteq \Sigma^*$ . Prove or refute:

- $(L_1 \cup L_2)^\omega \subseteq L_1^\omega \cup L_2^\omega$
  - $(L_1 \cup L_2)^\omega \supseteq L_1^\omega \cup L_2^\omega$
- $\lim(L_1 \cup L_2) \subseteq \lim L_1 \cup \lim L_2$
  - $\lim(L_1 \cup L_2) \supseteq \lim L_1 \cup \lim L_2$
- $L^\omega \subseteq \lim L^+$
  - $L^\omega \supseteq \lim L^+$
- $\lim(L_1 \cdot L_2) \subseteq L_1 \cdot L_2^\omega$
  - $\lim(L_1 \cdot L_2) \supseteq L_1 \cdot L_2^\omega$