



Automata and Logic

Exercise Sheet 3

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Exercise 12

Let $\Sigma := \{a, b\}$, $M := \{0, 1, 2\}$, and let $\circ: M \times M \rightarrow M$ be defined as $x \circ y := (x + y) \bmod 3$. We define mappings $\Phi, \Phi': \Sigma^* \rightarrow M$ by setting $\Phi(w) := |w| \bmod 3$ and $\Phi'(w) := |w|_a \bmod 3$, where $|w|$ denotes the *length* of w and $|w|_a$ the number of occurrences of the symbol a in w .

- Show that both Φ and Φ' are monoid homomorphisms from $(\Sigma^*, \cdot, \varepsilon)$ into $(M, \circ, 0)$.
- For each of the languages $\Phi^{-1}(\{0, 2\})$, $\Phi^{-1}(\{1\})$ and $(\Phi')^{-1}(\{1\})$ devise a finite automaton that recognises the language.

Exercise 13

Let Σ be an alphabet, $L \subseteq \Sigma^*$ a language, and $(M, \circ, 1)$ a monoid. Prove that L is accepted by $(M, \circ, 1)$, if and only if \bar{L} is also accepted by $(M, \circ, 1)$.

Exercise 14

Determine the syntactic monoid of the language described by a^*ba^* .

Exercise 15

Let $L \subseteq \Sigma^*$, and \approx be an equivalence relation on Σ^* . Consider the following property:

For all $u, v \in \Sigma^*$, if $u \in L$ and $u \approx v$, then $v \in L$. (*)

- The proof of Corollary 1.13 from the lecture depends on the fact that the syntactical congruence \sim_L has property (*). Prove this.
- Show that \sim_L is the coarsest congruence relation with property (*).
- Show that the Nerode right congruence ρ_L is the coarsest right congruence with property (*).

Note: An equivalence relation \approx_2 is *coarser* than \approx_1 if for every x, y , $x \approx_1 y$ implies $x \approx_2 y$. (In particular, \approx_2 has at most as many equivalence classes as \approx_1 .)

Exercise 16

Show that any submonoid of a finite group is also a group.

Exercise 17

Let V be an M -variety. Show that $L(V)_\Sigma$ is closed under union *without* using Thm. 1.22 from the lecture.

Exercise 18

Let Σ be an alphabet. Prove or refute the following claims:

- Every regular language $L \subseteq \Sigma^*$ is accepted by its syntactic monoid.
- If $L \subseteq \Sigma^*$ is accepted by a finite group, then the syntactic monoid of L is a finite group.
- For every regular language $L \subseteq \Sigma^*$, the syntactic monoid M_L is the smallest monoid accepting L ; i.e. for every monoid M accepting L , we have $|M_L| \leq |M|$.
- For a word $w = a_1 \dots a_n$, let \overleftarrow{w} denote the mirror image of w , i.e. $\overleftarrow{w} = a_n \dots a_1$. For a language $L \subseteq \Sigma^*$, we define $\overleftarrow{L} := \{\overleftarrow{w} \mid w \in L\}$. **Claim:** If the minimal automaton for L has n states, then the minimal automaton for \overleftarrow{L} has also n states.

Exercise 19

Let L_1 be the language over $\{a\}$ described by a^+ , and let L_2 be the language over $\{a, b\}$ described by $(a + b)^* b (a + b)^*$.

- Is there a monoid that accepts both L_1 and L_2 ?
- Are the syntactic monoids of those languages isomorphic?

Exercise 20

Let L_1 and L_2 be two languages over the same alphabet Σ that are accepted by the same monoid $(M, \circ, 1)$. Prove or refute the following statements:

- M accepts $L_1 \cap L_2$.
- M accepts $L_1 \cup L_2$.
- M accepts $L_1 \cdot L_2$.