

Faculty of Computer Science Institute of Theoretical Computer Science, Chair of Automata Theory

Automata and Logic

Exercise Sheet 6

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Exercise 35

Let $\Sigma := \{a\}$. Recall from the lecture that $L_{k,n}$ denotes the set of all first-order formulae (over the non-logical symbols =, <, and Q_a) containing k free variables and having quantifier depth at most n. For the following combinations of k, n, determine a *finite* set $\Gamma_{k,n}$ such that for every formula $\phi \in L_{k,n}$, there is a formula $\psi \in \Gamma_{k,n}$ with $\phi \equiv \psi$. Determine also the equivalence classes of $\equiv_{k,n}$.

- a) k = 1, n = 0;
- b) k = 2, n = 0;
- c) k = 0, n = 1; and
- d) k = 1, n = 1.

Exercise 36

Give the formulae ϕ_W for each equivalence class W of $\equiv_{2,0}$. Then determine a finite disjunction of formulae ϕ_W for $\equiv_{2,0}$ -classes which is equivalent to the formulae:

- a) true;
- b) $\neg(x < y) \lor x = y$; and
- c) false.

Exercise 37

Consider the Ehrenfeucht-Fraïssé games on the words

- a) ab and ba; and
- b) aaabaaa and aabaaa.

Determine the $k \in \{1, ..., 4\}$ such that Player I has a winning strategy in k moves.

Exercise 38

Consider the Ehrenfeucht-Fraïssé games on the words a^i and a^j with i < j.

- a) Describe an optimal winning strategy for Player I, i.e., a strategy such that Player I wins with a minimal number of moves.
- b) Prove that Player I has a winning strategy on a^i and a^j (with i < j) in *m* moves if $i < 2^m 1$.