Lehrstuhl für Automatentheorie
Technische Universität Dresden
Prof. Dr. F. Baader

01187 Dresden
Tel.: 0351/463–38314

15. Exercises for the lecture **Term Rewriting Systems**

Exercise 62:

The semi-decision procedure described in the proof of Theorem 7.7 of the lecture is rather inefficient: for the input $s \approx_E t$, all R_i-normal forms of s and t are computerd in the i'th iteration of the "Repeat"-loop. Show that the following modification of the procedure still yields a partial decision procedure for the word problem:

- Begin with $s_0 := s$ and $t_0 := t$.
- After the ith "Repeat"-loop, compute *one arbitrary* R_i-normal form s_i of s_{i_1} and *one arbitrary* R_i-normal form t_i of t_{i_1}.
- Output "yes"($s \approx_E t$) iff there exists an n such that $s_n \approx t_n$.

Hint: show: (1) for all $n \geq 0$ it holds that $s \approx_E s_n$ and $t \approx_E t_n$; (2) since R_∞ is terminating, there exists an n such that $s_n = s_m$ and $t_n = t_m$ for all $m \geq n$.

Exercise 63:

(a) Consider the following set of identities:

$$E := \{f(f(x,y),z) \approx f(x,f(y,z)),\ f(x,x) \approx x,\ f(f(x,y),x) \approx x\}$$

Apply the rules of the improved completion procedure to E. Use a strategy that resembles the basic completion procedure, but simplifies rules as follows: upon adding new rules, simplify the old ones by means of L-Simplify-Rule and R-Simplify-Rule.

(b) Consider the following set of identities:

$$E := \{x + (y + z) \approx (x + y) + z,\ f(x) + f(y) \approx f(x + y)\}$$

Apply the completion procedure described above to the input E and the polynomial order induced by

$$P_f(X) = X + 1,\ P_+(X,Y) = XY^2$$

Exercise 64:

Consider the following completion procedure for ground term rewriting systems (and also recall Exercise 54):

Input	G_0, a finite set of ground identities over Σ, $>$, a reduction order total on the set of ground terms over Σ.
Procedure	Apply the rules L-Simplify-Rule, Delete and Orient, until no rule is applicable any more.
Output	a ground term rewriting system.

Show that this procedure:
(a) always terminates,
(b) is fair,
(c) is correct and
(d) never ends in failure.