

Chapter 4

Reasoning with Tableaux Algorithms

Outline

We start with an algorithm for deciding **consistency of an ABox without a TBox**.

This covers most of the inference problems introduced in Chapter 2:

- acyclic TBoxes can be eliminated by **expansion**
- satisfiability, subsumption, equivalence, and instance checking can be reduced to ABox consistency

Idea

The **tableau-based consistency algorithm** tries to generate a **finite model** for the input ABox \mathcal{A}_0

- **tableau rules** extend the ABox (one rule **per constructor**)
- **obvious contradictions** show inconsistency
- if an ABox is:
complete (no rule applicable) and
open (no contradictions found),
then it describes a model

Example

\mathcal{T} : Heroine \equiv Hero \sqcap Female

Subsumption:

$\exists \text{helps.Hero} \sqcap \exists \text{helps.Female} \sqsubseteq_{\mathcal{T}}^? \exists \text{helps.Heroine}$

Reduction to satisfiability: is the following concept unsatisfiable w.r.t. \mathcal{T} ?

$\exists \text{helps.Hero} \sqcap \exists \text{helps.Female} \sqcap \neg \exists \text{helps.Heroine}$

Reduction to consistency: is the following ABox inconsistent w.r.t. \mathcal{T} ?

$\{(\exists \text{helps.Hero} \sqcap \exists \text{helps.Female} \sqcap \neg \exists \text{helps.Heroine})(a)\}$

Expansion: is the following ABox inconsistent?

$\{(\exists \text{helps.Hero} \sqcap \exists \text{helps.Female} \sqcap \neg \exists \text{helps.}(\text{Hero} \sqcap \text{Female}))(a)\}$

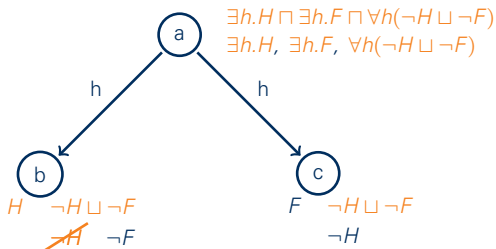
NNF: is the following ABox inconsistent?

$\{(\exists \text{helps.Hero} \sqcap \exists \text{helps.Female} \sqcap \forall \text{helps.}(\neg \text{Hero} \sqcup \neg \text{Female}))(a)\}$

Expansion of the ABox

Deciding inconsistency of

$$\{(\exists \text{helps.Hero} \sqcap \exists \text{helps.Female} \sqcap \forall \text{helps.}(\neg \text{Hero} \sqcup \neg \text{Female}))(\text{a})\}$$



This is a **complete, open** ABox \rightarrow **model** of input ABox

It is **consistent**, subsumption **does not hold**

Formal Algorithm

Input An ABox \mathcal{A}_0

Output "yes" if \mathcal{A}_0 is consistent, "no" otherwise

Preprocessing: transform all concept descriptions of \mathcal{A}_0 to **negation normal form**

$$\neg(C \sqcap D) \rightsquigarrow \neg C \sqcup \neg D$$

$$\neg(C \sqcup D) \rightsquigarrow \neg C \sqcap \neg D$$

$$\neg\neg C \rightsquigarrow C$$

$$\neg(\exists r.C) \rightsquigarrow \forall r.\neg C$$

$$\neg(\forall r.C) \rightsquigarrow \exists r.\neg C$$

NNF transformation in **polynomial time**, is **semantics invariant**

Formal Algorithm (2)

Data Structure:

finite set of ABoxes. Initialized to $\{\mathcal{A}_0\}$ (in NNF)

Rule applications:

tableau rules replace one ABox from the set by finitely many new ABoxes

Termination:

when **no** rule can be applied to any ABox in the set

ABox is **complete** if no rule applies to it

Return:

“**yes**” if the set contains an **open** ABox

\mathcal{A} is open if contains no obvious contradiction of the form $A(a), \neg A(a)$

“**no**” otherwise (i.e., if all ABoxes are **closed**)

Tableau Rules

There is one rule for each constructor (except negation)

\sqcap -rule

Condition \mathcal{A} contains $(C \sqcap D)(a)$ but not both $C(a)$ and $D(a)$

Action $\mathcal{A}' := \mathcal{A} \cup \{C(a), D(a)\}$

\sqcup -rule

Condition \mathcal{A} contains $(C \sqcup D)(a)$ but neither $C(a)$ nor $D(a)$

Action $\mathcal{A}' := \mathcal{A} \cup \{C(a)\}$ and $\mathcal{A}'' := \mathcal{A} \cup \{D(a)\}$

\exists -rule

Condition \mathcal{A} contains $(\exists r.C)(a)$ but there is no b with $\{r(a, b), C(b)\} \subseteq \mathcal{A}$

Action $\mathcal{A}' := \mathcal{A} \cup \{r(a, b), C(b)\}$ where b is a **new** individual name

\forall -rule

Condition \mathcal{A} contains $(\forall r.C)(a)$ and $r(a, b)$ but not $C(b)$

Action $\mathcal{A}' := \mathcal{A} \cup \{C(b)\}$

Tableau Algorithm is a Decision Procedure

Lemma 4.1
rules preserve consistency

Lemma 4.10
termination

Lemma 4.4

soundness complete and open = consistent
completeness closed = inconsistent

