

Faculty of Computer Science Institute of Theoretical Computer Science, Chair of Automata Theory

# **Description Logics**

### **Exercise Sheet 8**

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## Exercise 31

Let  $\mathcal{K} = \langle \mathcal{A}_0, \mathcal{T} \rangle$  be an  $\mathcal{ALC}$ -knowledge base, where  $\mathcal{T}$  is a general TBox. The *precompletion* of  $\mathcal{K}$  is the set of ABoxes M that is produced by the tableau algorithm when starting with the set of ABoxes  $\{\mathcal{A}_0\}$  and exhaustively applying all tableau rules plus the GCI-rule except for the modified  $\exists$ -rule. Do the following:

a) Show that  $\mathcal{K}$  is consistent iff there is an open ABox  $\mathcal{A} \in M$  such that for all individual names *a* occurring in  $\mathcal{A}$ , the concept  $C^a_{\mathcal{A}} := \prod_{C(a) \in \mathcal{A}} C$  is satisfiable w.r.t.  $\mathcal{T}$ .

**Hint:** For the "if" direction, proceed as follows: The correctness of the tableau algorithm for  $\mathcal{ALC}$  implies that, if  $C^a_{\mathcal{A}}$  is satisfiable, then exhaustively applying all (!) rules to the set of ABoxes {{ $C^a_{\mathcal{A}}(a)$ } yields a set M' that contains an open and complete ABox. Show how to join all these ABoxes to obtain an open and complete tableau for  $\mathcal{A}$  and conclude that  $\mathcal{A}_0$  is consistent w.r.t.  $\mathcal{T}$ .

b) Use the result from a) to prove that ABox consistency in *ALC* can be decided in deterministic exponential time.

# Exercise 32

For each of the following languages of binary trees over the alphabet  $\Sigma = \{a, b\}$ , define a looping tree automaton that accepts the language.

- a) The set of all trees that contain a branch (starting at the root) in which all nodes are labelled with *a*.
- b) The set of all trees that do not contain nodes  $n_0$ ,  $n_1$ ,  $n_2$  such that
  - $n_1 = n_0 i$  for some  $i \in \{0, 1\}$ ,
  - $n_2 = n_1 j$  for some  $j \in \{0, 1\}$ , and
  - $T(n_0) = T(n_1) = T(n_2) = a$ .

### Exercise 33

Show that there is no looping tree automaton on binary  $\{a, b\}$ -trees that accepts the set of all trees that contain a branch with infinitely many nodes labelled with a.