



Description Logics

Exercise Sheet 10

Dr.-Ing. habil. Anni-Yasmin Turhan/Francesco Kriegel
Winter Semester 2014/15

Exercise 38

Consider the \mathcal{EL} -TBox \mathcal{T} consisting of the following axioms:

$$\begin{aligned} A &\sqsubseteq B \sqcap \exists r.C \\ B \sqcap \exists r.B &\sqsubseteq C \sqcap D \\ C &\sqsubseteq \exists r.A \sqcap B \\ \exists r.\exists r.B \sqcap D &\sqsubseteq \exists r.(A \sqcap B) \end{aligned}$$

Check whether the following subsumption relations hold w.r.t. \mathcal{T} :

- a) $A \sqsubseteq B$
- b) $A \sqsubseteq \exists r.\exists r.A$
- c) $B \sqcap \exists r.A \sqsubseteq \exists r.C$

Exercise 39

Determine whether Player 2 has a winning strategy in the PSPACE game $G = (\phi, \{p_0, p_2\}, \{p_1, p_3\}, <)$ with

$$\phi = (\neg p_0 \rightarrow p_1) \wedge ((p_0 \wedge p_1) \rightarrow (p_2 \vee p_3)) \wedge (\neg p_1 \rightarrow (p_3 \rightarrow \neg p_2))$$

and $p_i < p_j$ iff $i < j$.

Exercise 40

A *quantified Boolean formula* is of the form $\phi = Q_1 p_1 \dots Q_n p_n . \psi$ where $Q_1, \dots, Q_n \in \{\forall, \exists\}$ are quantifiers, p_1, \dots, p_n are propositional variables, and ψ is a propositional formula containing only the variables p_1, \dots, p_n .

Validity of such formulae is defined as follows:

- For $n = 0$, the formula ϕ does not contain variables, and thus is a Boolean combination of 0 and 1. Then ϕ is valid iff it evaluates to 1.
- For $n > 0$, we consider:

$$\begin{aligned} \phi_0 &:= Q_2 p_2 \dots Q_n p_n . \psi[p_1 := 0], \text{ and} \\ \phi_1 &:= Q_2 p_2 \dots Q_n p_n . \psi[p_1 := 1]. \end{aligned}$$

If $Q_1 = \exists$, then ϕ is valid iff one of ϕ_0 and ϕ_1 is valid.

If $Q_1 = \forall$, then ϕ is valid iff both ϕ_0 and ϕ_1 are valid.

QBF denotes the set of all valid quantified Boolean formulae. Prove by reduction of QBF that the problem of deciding the existence of winning strategy for Player 2 in PSPACE games is PSPACE-hard.

Exercise 41

Finish the proof of Lemma 6.7 by showing that $v_0 \in C_G^I$.