



Description Logics

Exercise Sheet 10

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Exercise 37

Let \mathcal{T} be a general \mathcal{EL} -TBox and $\hat{\mathcal{T}}$ the TBox obtained from exhaustive application of the normalisation rules NF1–NF5 to \mathcal{T} .

Finish the proof of Lemma 5.18 by showing the following:

- $\hat{\mathcal{T}}$ can be obtained from \mathcal{T} in polynomial time.
- $\hat{\mathcal{T}}$ is in normal form.
- For all concept names A, B occurring in \mathcal{T} , we have $A \sqsubseteq_{\mathcal{T}} B$ iff $A \sqsubseteq_{\hat{\mathcal{T}}} B$.

Exercise 38

Consider the \mathcal{EL} -TBox \mathcal{T} consisting of the following axioms:

$$\begin{aligned} A &\sqsubseteq B \sqcap \exists r.C \\ B \sqcap \exists r.B &\sqsubseteq C \sqcap D \\ C &\sqsubseteq \exists r.A \sqcap B \\ \exists r.\exists r.B \sqcap D &\sqsubseteq \exists r.(A \sqcap B) \end{aligned}$$

Check whether the following subsumption relations hold w.r.t. \mathcal{T} :

- $A \sqsubseteq B$
- $A \sqsubseteq \exists r.\exists r.A$
- $B \sqcap \exists r.A \sqsubseteq \exists r.C$

Exercise 39

A *role complement* is a role of the form $\neg r$, where r is a role name. The semantics of role complements is defined as follows:

$$(\neg r)^{\mathcal{I}} := \Delta^{\mathcal{I}} \times \Delta^{\mathcal{I}} \setminus r^{\mathcal{I}}.$$

The description logic \mathcal{ALC}^{\neg} extends \mathcal{ALC} by role complements, i.e. role complements are allowed to occur in existential restrictions, value restrictions and role assertions.

Show that \mathcal{ALC}^{\neg} does not have the tree-model property.

Exercise 40

Use a tableau algorithm to decide whether the following \mathcal{ALC} -knowledge base is consistent:

$$\begin{aligned}\mathcal{T} &:= \{A \sqcap \forall r. \neg A \sqsubseteq \perp\} \\ \mathcal{A} &:= \{(\forall r. \neg A)(a), (\exists r. A)(b), r(a, b)\}\end{aligned}$$

Exercise 41

For each of the following \mathcal{ALC} -concept descriptions C and \mathcal{ALC} -TBoxes \mathcal{T} decide whether C is satisfiable w.r.t. \mathcal{T} by constructing the looping tree automaton $\mathcal{A}_{C, \mathcal{T}}$ and checking its accepted language $L(\mathcal{A}_{C, \mathcal{T}})$ for emptiness.

- a) $C := A$
 $\mathcal{T} := \{A \sqsubseteq \neg A\}$
- b) $C := A$
 $\mathcal{T} := \emptyset$
- c) $C := A \sqcap \exists r. A$
 $\mathcal{T} := \{A \sqsubseteq \forall r. \neg A\}$