



Description Logic

Winter Semester 2017/18

Exercise Sheet 12

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Exercise 12.1 We consider simulations, which are “one-sided” variants of bisimulations. Given interpretations \mathcal{I}_1 and \mathcal{I}_2 , the relation $\sigma \subseteq \Delta^{\mathcal{I}_1} \times \Delta^{\mathcal{I}_2}$ is a *simulation* between \mathcal{I}_1 and \mathcal{I}_2 if

- whenever $d_1 \sigma d_2$ and $d_1 \in A^{\mathcal{I}_1}$, then $d_2 \in A^{\mathcal{I}_2}$, for all $d_1 \in \Delta^{\mathcal{I}_1}$, $d_2 \in \Delta^{\mathcal{I}_2}$, and $A \in \mathbf{C}$;
- whenever $d_1 \sigma d_2$ and $(d_1, d'_1) \in r^{\mathcal{I}_1}$, then there exists a $d'_2 \in \Delta^{\mathcal{I}_2}$ such that $d'_1 \sigma d'_2$ and $(d_2, d'_2) \in r^{\mathcal{I}_2}$, for all $d_1, d'_1 \in \Delta^{\mathcal{I}_1}$, $d_2 \in \Delta^{\mathcal{I}_2}$, and $r \in \mathbf{R}$.

We write $(\mathcal{I}_1, d_1) \approx (\mathcal{I}_2, d_2)$ if there is a simulation σ between \mathcal{I}_1 and \mathcal{I}_2 such that $d_1 \sigma d_2$.

- Show that $(\mathcal{I}_1, d_1) \sim (\mathcal{I}_2, d_2)$ implies $(\mathcal{I}_1, d_1) \approx (\mathcal{I}_2, d_2)$ and $(\mathcal{I}_2, d_2) \approx (\mathcal{I}_1, d_1)$.
- Is the converse of the implication in (a) also true?
- Show that, if $(\mathcal{I}_1, d_1) \approx (\mathcal{I}_2, d_2)$, then for all \mathcal{EL} -concepts C it holds that $d_1 \in C^{\mathcal{I}_1}$ implies $d_2 \in C^{\mathcal{I}_2}$.
- Which of the constructors disjunction, negation, or value restriction can be added to \mathcal{EL} without losing the property in (c)?
- Show that \mathcal{ALC} is more expressive than \mathcal{EL} .
- Show that \mathcal{ELI} is more expressive than \mathcal{EL} .
- Can the fact that subsumption in \mathcal{EL} is decidable in polynomial time, while subsumption in \mathcal{ELI} is EXPTIME-complete, be used to show that \mathcal{ELI} is more expressive than \mathcal{EL} ?

Exercise 12.2 Consider the TBox

$$\mathcal{T} = \{A_1 \sqcap A_2 \sqsubseteq \exists r.B, \exists r^-.A_2 \sqsubseteq C, A \sqsubseteq A_1 \sqcap A_2, \exists r.(B \sqcap C) \sqsubseteq D\},$$

where A, A_1, A_2, B, C, D are concept names. Use the classification procedure for \mathcal{ELI} to check whether the following subsumption relationships hold w.r.t. \mathcal{T} :

- $A \sqsubseteq D$
- $\exists r.A \sqsubseteq \exists r.D$
- $A \sqsubseteq \exists r.A$