

Query Rewriting for *DL-Lite* with n -ary Concrete Domains

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TECHNISCHE
UNIVERSITÄT
DRESDEN

2

TNG



TECHNOLOGY
CONSULTING



Measurements

Patient	Heart rate	Age	Timestamp
p78435	80	45	2017-06-09T11:46:50
p71247	60	60	2017-06-09T11:47:05
p72349	120	85	2017-06-09T11:48:19
p73872	160	80	2017-06-09T12:23:16
⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮
⋮	⋮	⋮	⋮



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Select Patient From Measurements Where **HR > 140**



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$$\exists h. \text{Patient}(p) \wedge \text{HR}(p, h) \wedge \mathbf{h > 140}$$



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Select Patient From Measurements Where **HR > 140**

$$\exists h. \text{Patient}(p) \wedge \text{HR}(p, h) \wedge \mathbf{h > 140}$$

$$\text{MaxHR} = 220 - \text{Age} \quad \text{HR} = \text{MaxHR} \rightarrow \text{Critical!}$$



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Patient(p) $\wedge$ **Critical(p)**

MaxHR = $220 - \text{Age}$ HR = MaxHR $\rightarrow$ Critical!



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Select Patient From Measurements Where **HR = 220 - Age**

Patient(p) $\wedge$ **Critical(p)**

rewritable?

MaxHR = 220 - Age

HR = MaxHR $\rightarrow$ Critical!

DL-Lite with n -ary Concrete Domains

Concrete domain $\mathcal{D}$

domain $\Delta^{\mathcal{D}}$

predicates $P_i^{\mathcal{D}} \subseteq (\Delta^{\mathcal{D}})^{n_i}$

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$$(\mathbf{P}_{220-})^{\mathcal{D}} := \{(x, y) \mid x = 220 - y\}$$

$$(\mathbf{Eq})^{\mathcal{D}} := \{(x, y) \mid x = y\}$$

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$DL\text{-}Lite_{horn}^{(\mathcal{HF})}(\mathcal{D})$

Patient, Doctor, Critical, ...

(atomic concepts)

HasDoctor, ...

(roles)

p73872, p78435, ...

(individuals)

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p73872, p78435, ...

HR, Age, ...

160, 220, ...

(atomic concepts)

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160, 220, ...

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ABox $\mathcal{A}$: Patient(p73872) HR(p73872, 160) Age(p73872, 80)

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TBox $\mathcal{T}$: Patient $\sqsubseteq \exists \text{MaxHR}$ Patient $\sqsubseteq \forall \text{HR}, \text{HR.Eq}$

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Patient $\sqsubseteq \forall \text{MaxHR, Age.P}_{220-}$ $\exists \text{HR, MaxHR.Eq} \sqsubseteq \text{Critical}$

Concrete domain $\mathcal{D}$ ABox $\mathcal{A}$ TBox $\mathcal{T}$

Concrete domain $\mathcal{D}$ ABox $\mathcal{A}$ TBox $\mathcal{T}$ Query ϕ

$$\exists h, m. \text{Patient}(p) \wedge \text{HR}(p, h) \wedge \text{MaxHR}(p, m) \wedge \text{Eq}(h, m)$$



Concrete domain $\mathcal{D}$ ABox $\mathcal{A}$ TBox $\mathcal{T}$ Query ϕ

$\exists h, m. \text{Patient}(p) \wedge \text{HR}(p, h) \wedge \text{MaxHR}(p, m) \wedge h = m$



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$\text{cert}(\mathcal{A}, \mathcal{T}, \phi) = ?$

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Rewritability

Data complexity

first-order rewritable: $\text{cert}(\mathcal{A}, \mathcal{T}, \phi) = \text{ans}(\mathcal{A}, \phi_{\mathcal{T}})$

AC^0

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AC^0

combined rewritable: $\text{cert}(\mathcal{A}, \mathcal{T}, \phi) = \text{ans}(\mathcal{A}_{\mathcal{T}}, \phi_{\mathcal{T}})$

P

Conditions for Rewritability

Convex concrete domains

(Artale, Ryzhikov, and Kontchakov 2012; Baader, Brandt, and Lutz 2005; Savković and Calvanese 2012)

If $\phi_1 \wedge \dots \wedge \phi_m$ implies $\psi_1 \vee \dots \vee \psi_n$, then it implies some ψ_j .

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$$x = 220 - y \rightarrow x < 200 \vee y \leq 20 \quad \textcolor{red}{\text{!}}$$

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$$x = 220 - y \rightarrow x < 200 \vee y \leq 20 \nmid$$

Examples: $\mathbb{Q}$ with $\mathbf{Eq}$, $\mathbf{Eq}_q$, $\mathbf{P}_{q-}$
 $\mathbb{Q}$ with $\mathbf{Eq}$, $\mathbf{Eq}_q$, $\mathbf{P}_{q+}$, $>_q$ (or $<_q$)
 Σ^* with $\mathbf{Eq}$, $\mathbf{Eq}_w$, $\mathbf{concat}_w$

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cr-admissibility

(1) **equality** predicates $\mathbf{Eq}$, $\mathbf{Eq}_d$

$$x = y \quad x = 100$$

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 $\mathbb{Q}$ with Eq , Eq_q , P_{q+} , $>_q$ (or $<_q$)
 Σ^* with Eq , Eq_w , concat_w

cr-admissibility

- (1) equality predicates Eq , Eq_d $x = y$ $x = 100$
- (2) **decidable** in polynomial time $x = 220 - y \wedge y > 50 \rightarrow x < 100$?

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- (1) equality predicates $\mathbf{Eq}$, $\mathbf{Eq}_d$ $x = y$ $x = 100$
- (2) decidable in polynomial time $x = 220 - y \wedge y > 50 \rightarrow x < 100?$
- (3) **compute** counterexamples in polynomial time $x \mapsto 160, y \mapsto 60$

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- (3) compute counterexamples in polynomial time $x \mapsto 160, y \mapsto 60$
- (4) **functional** predicates $x = 220 - y \wedge y = 60 \rightarrow x = 160$

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- (1) equality predicates $\mathbf{Eq}$, $\mathbf{Eq}_d$ $x = y$ $x = 100$
- (2) decidable in polynomial time $x = 220 - y \wedge y > 50 \rightarrow x < 100?$
- (3) compute counterexamples in polynomial time $x \mapsto 160, y \mapsto 60$
- (4) functional predicates $x = 220 - y \wedge y = 60 \rightarrow x = 160$
- (5) **convexity**

The Rewriting

Implications in the TBox

$\text{Patient}(p) \wedge \text{Critical}(p)$

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 $\quad \quad \quad \leftarrow$

$\exists h, m. \text{Patient}(p) \wedge \text{HR}(p, h) \wedge \text{MaxHR}(p, m) \wedge h = m$

Implications in the TBox

$$\text{Patient}(p) \wedge \text{Critical}(p)$$

$$\exists \text{HR}, \text{MaxHR}. \text{Eq} \sqsubseteq \text{Critical}$$

$\leftarrow$

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Implications in the concrete domain

$$\exists m. \text{Patient}(p) \wedge \text{MaxHR}(p, m) \wedge m = 160$$

$$m = 220 - a \wedge a = 60 \rightarrow m = 160$$

Implications in the TBox

$$\text{Patient}(p) \wedge \text{Critical}(p)$$

$$\exists \text{HR}, \text{MaxHR}. \text{Eq} \sqsubseteq \text{Critical}$$

$\leftarrow$

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Implications in the concrete domain

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$\leftarrow$

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Implications in the TBox

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$\leftarrow$

$$\exists m, a. \text{Patient}(p) \wedge \text{MaxHR}(p, m) \wedge \text{Age}(p, a) \wedge m = 220 - a \wedge a = 60$$

Bound the number of new atoms!

Summary

Combined rewriting for cr-admissible concrete domains

First-order rewriting for unary predicates

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Open problems

First-order rewriting?

Polynomial query rewriting $\phi_{\mathcal{T}}$?

Query answering over concrete domains in other logics?

(Hernich, Lemos, and Wolter 2017)

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Thank you!

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